

Lecture 3

Binary Search Trees (contd.)

Inorder Tree Walk: Proof of Correctness

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Proof:

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Proof: **Basis Step:** Statement is trivially true when BST contains only **one** node.

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Inductive Hypothesis (IH)



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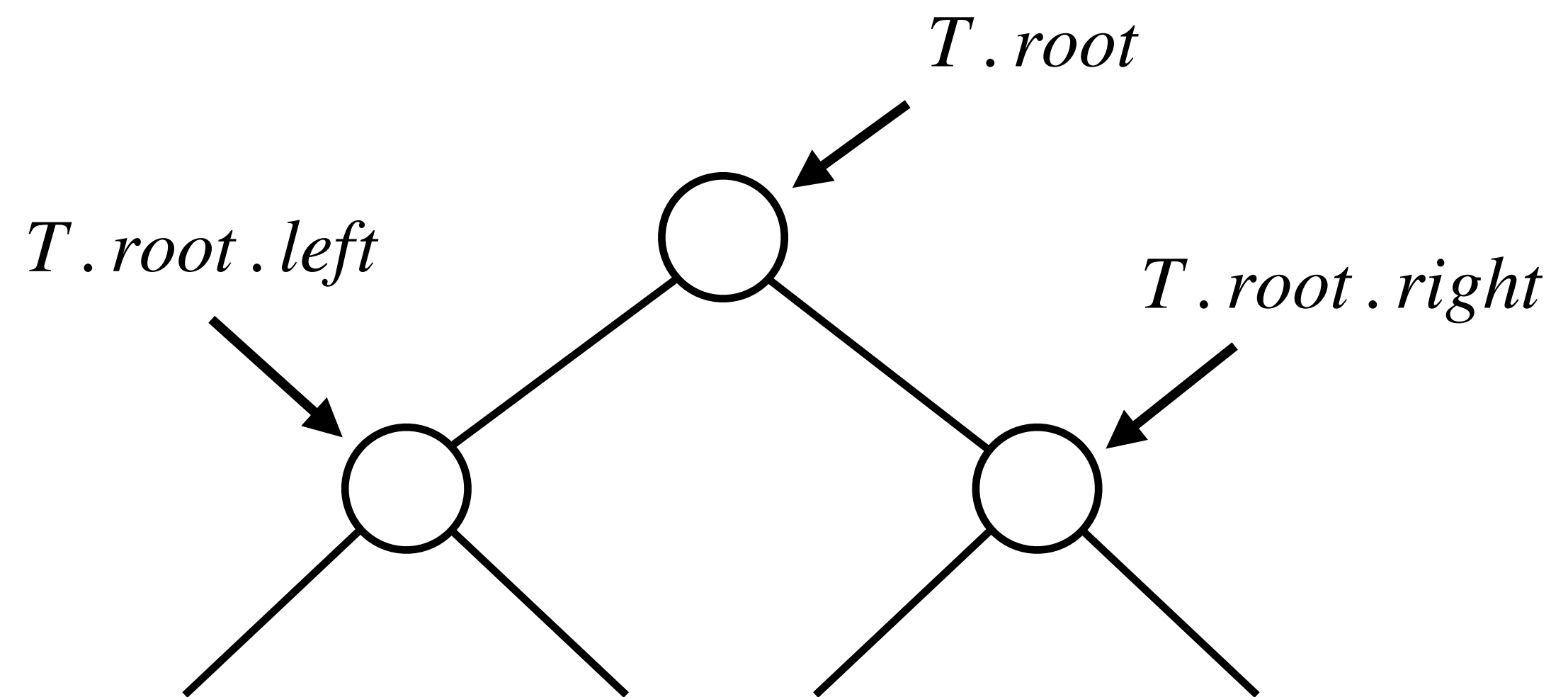
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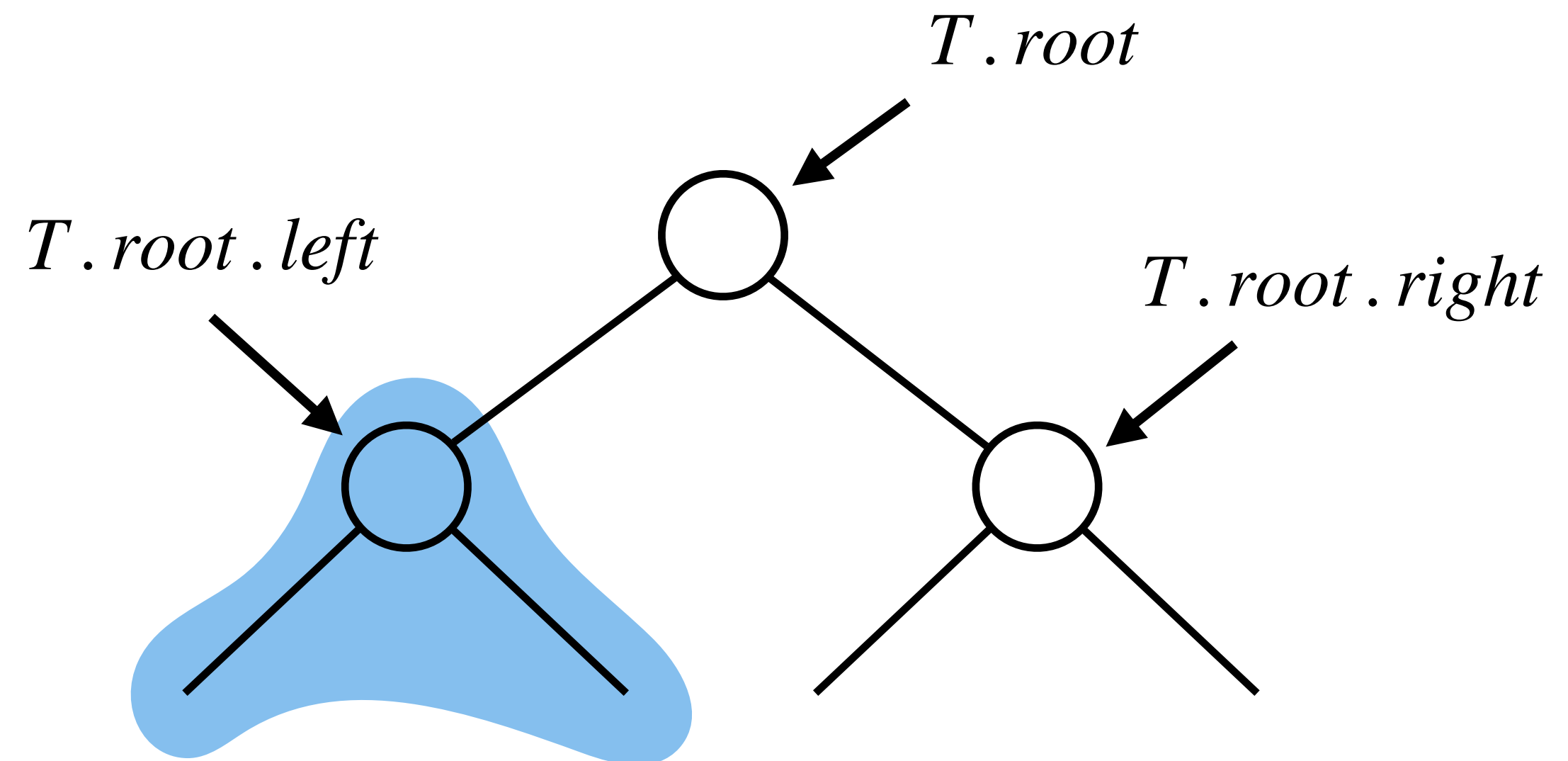
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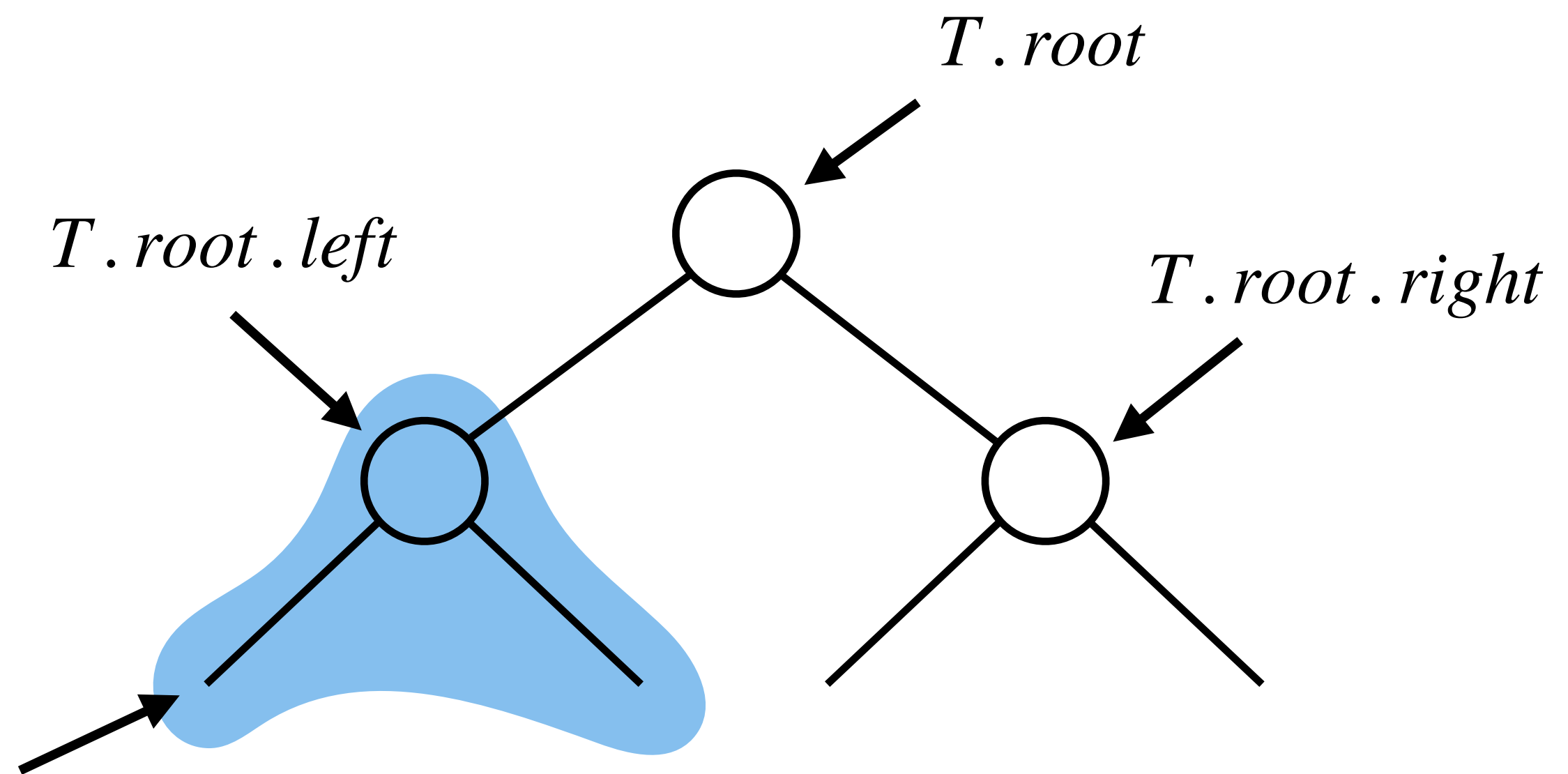
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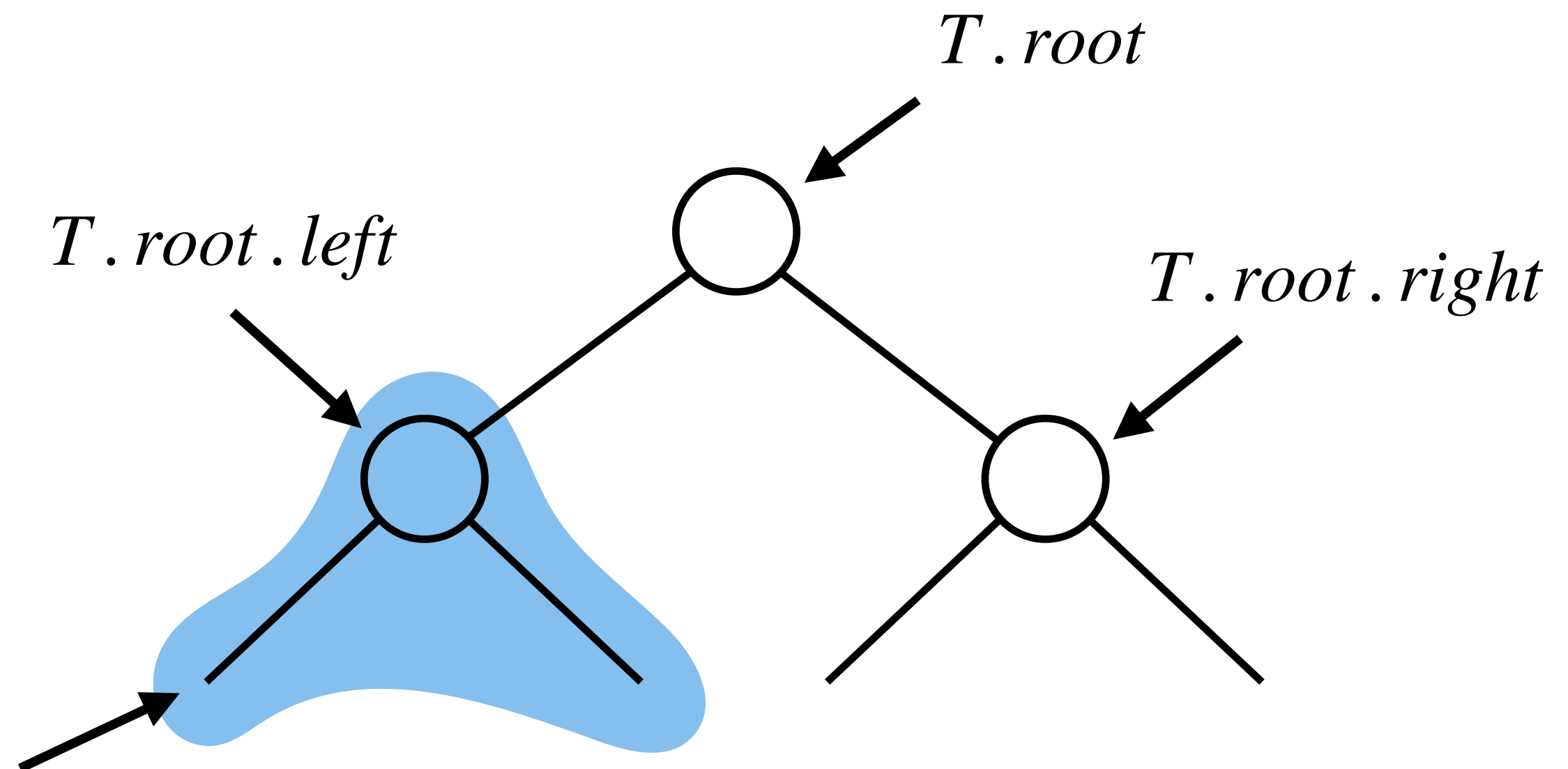
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First print keys here in sorted order
(\because IH and this subtree has $\leq k$ keys)



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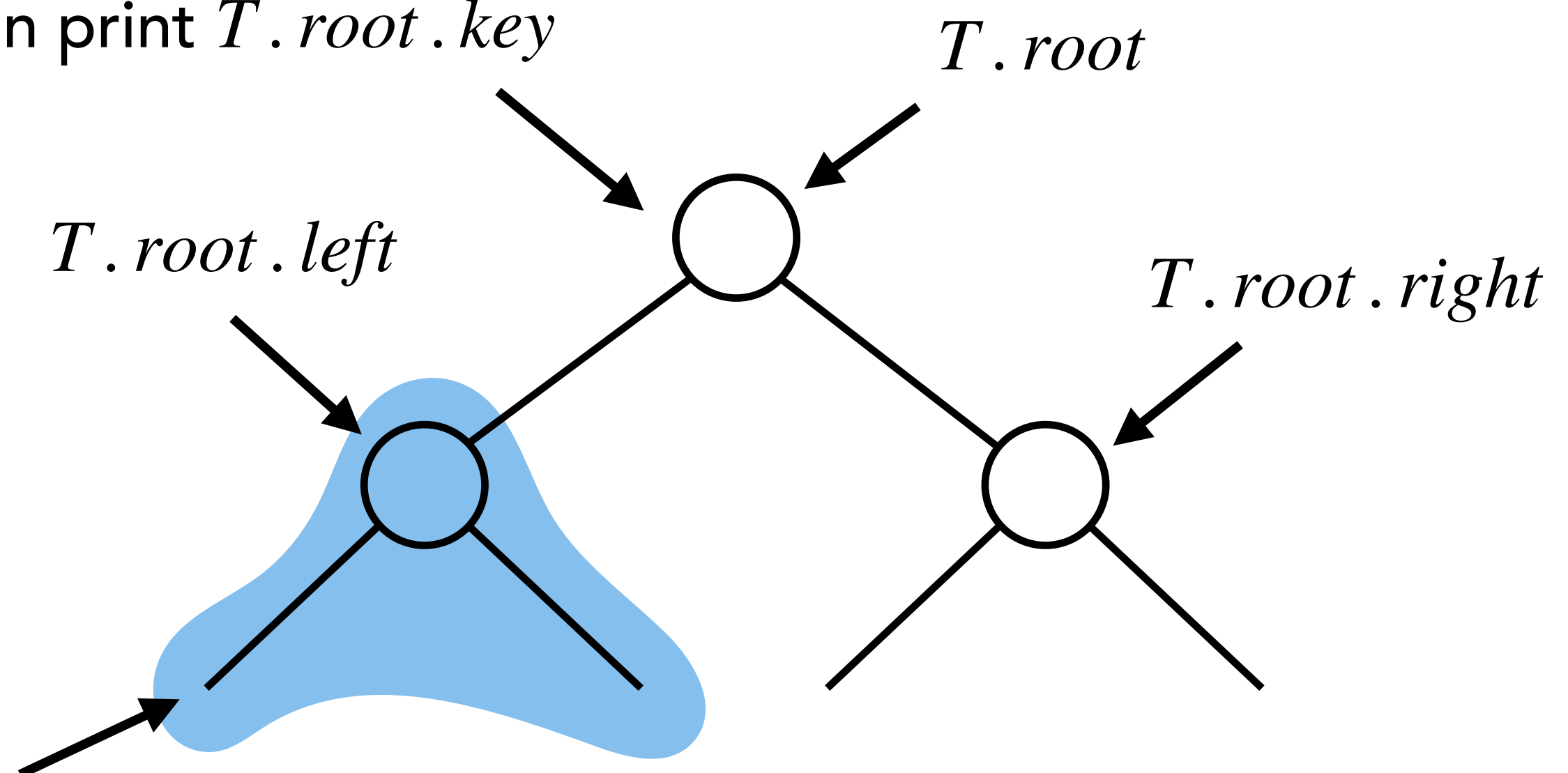
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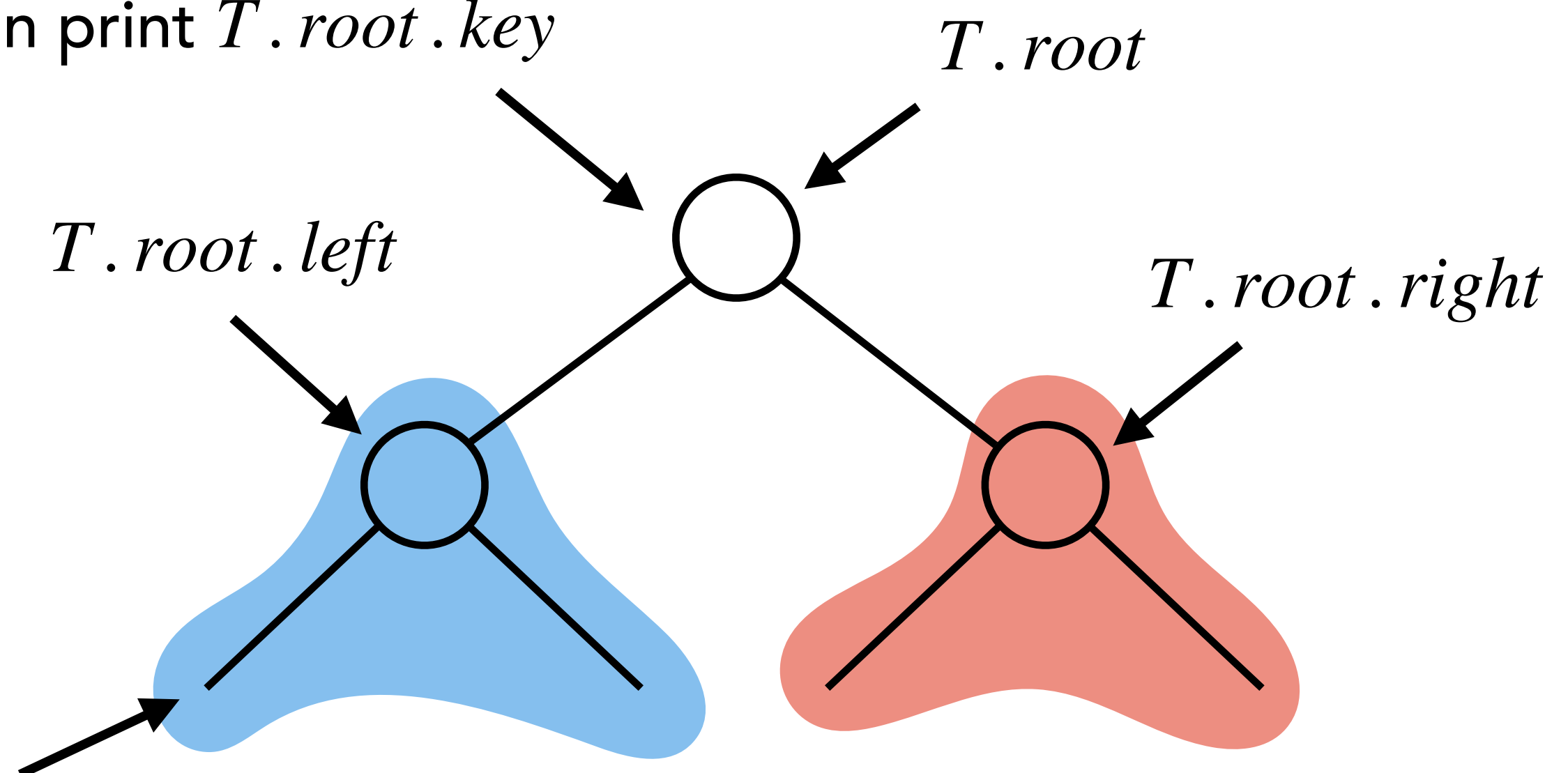
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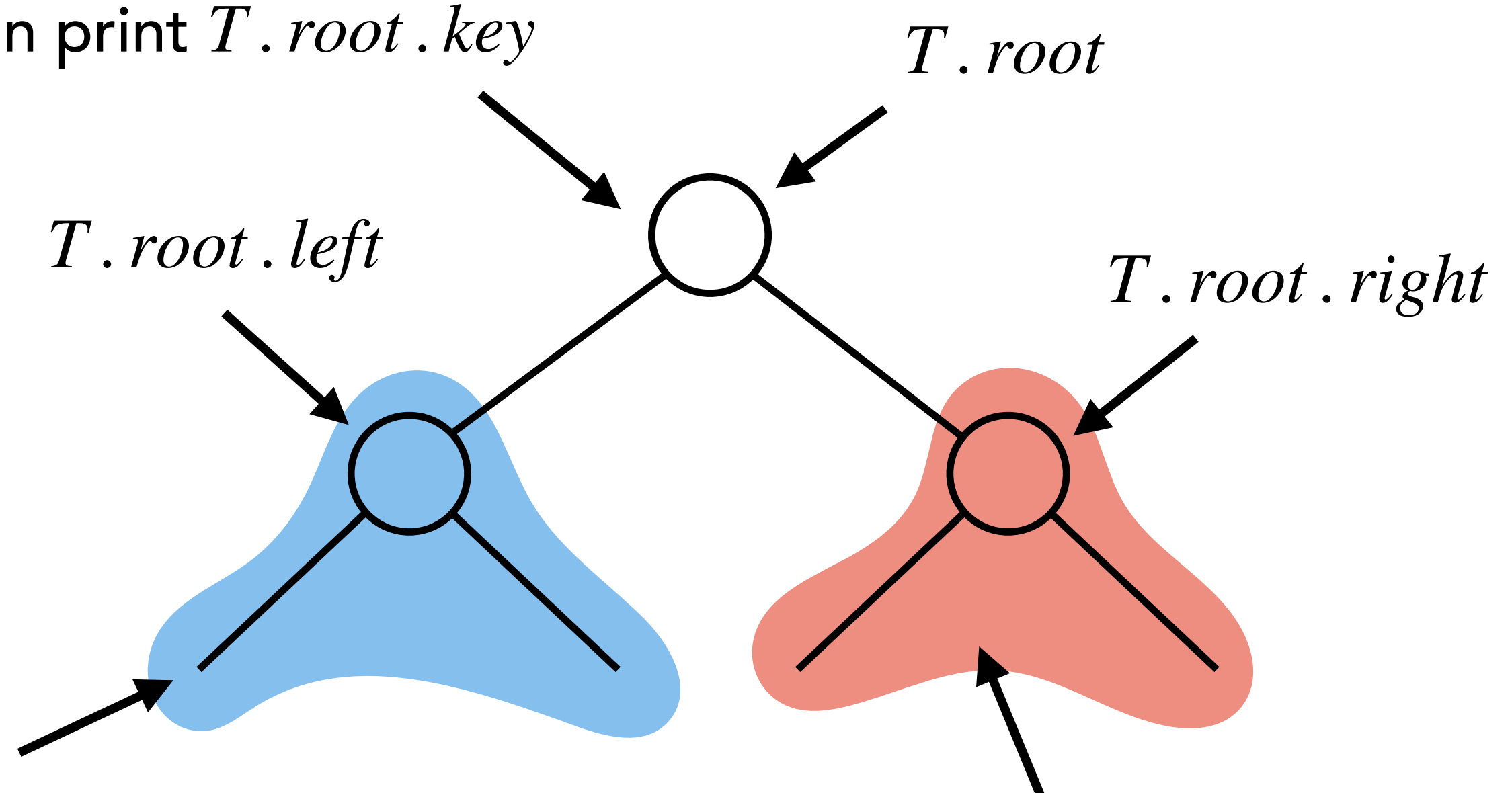
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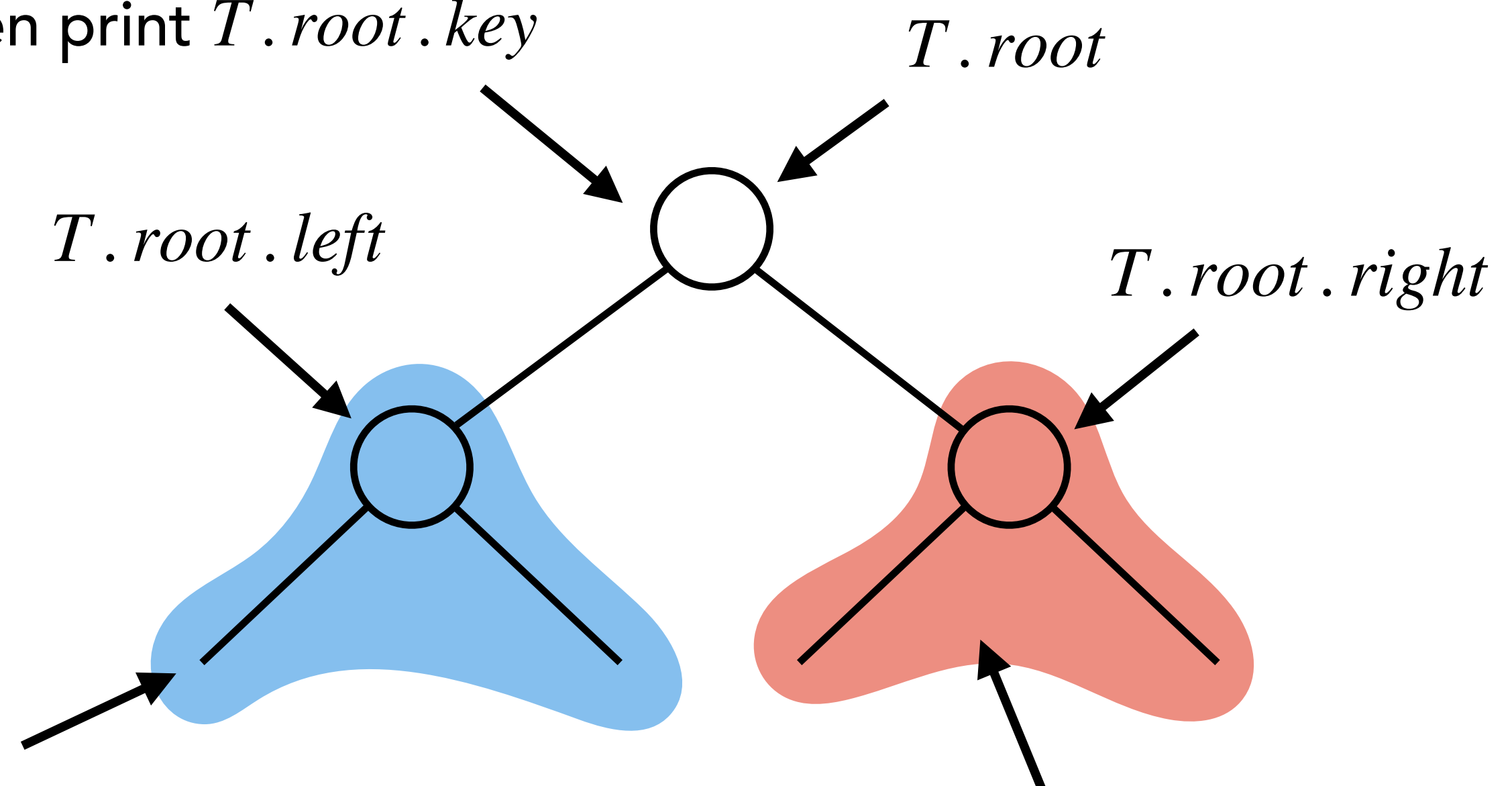
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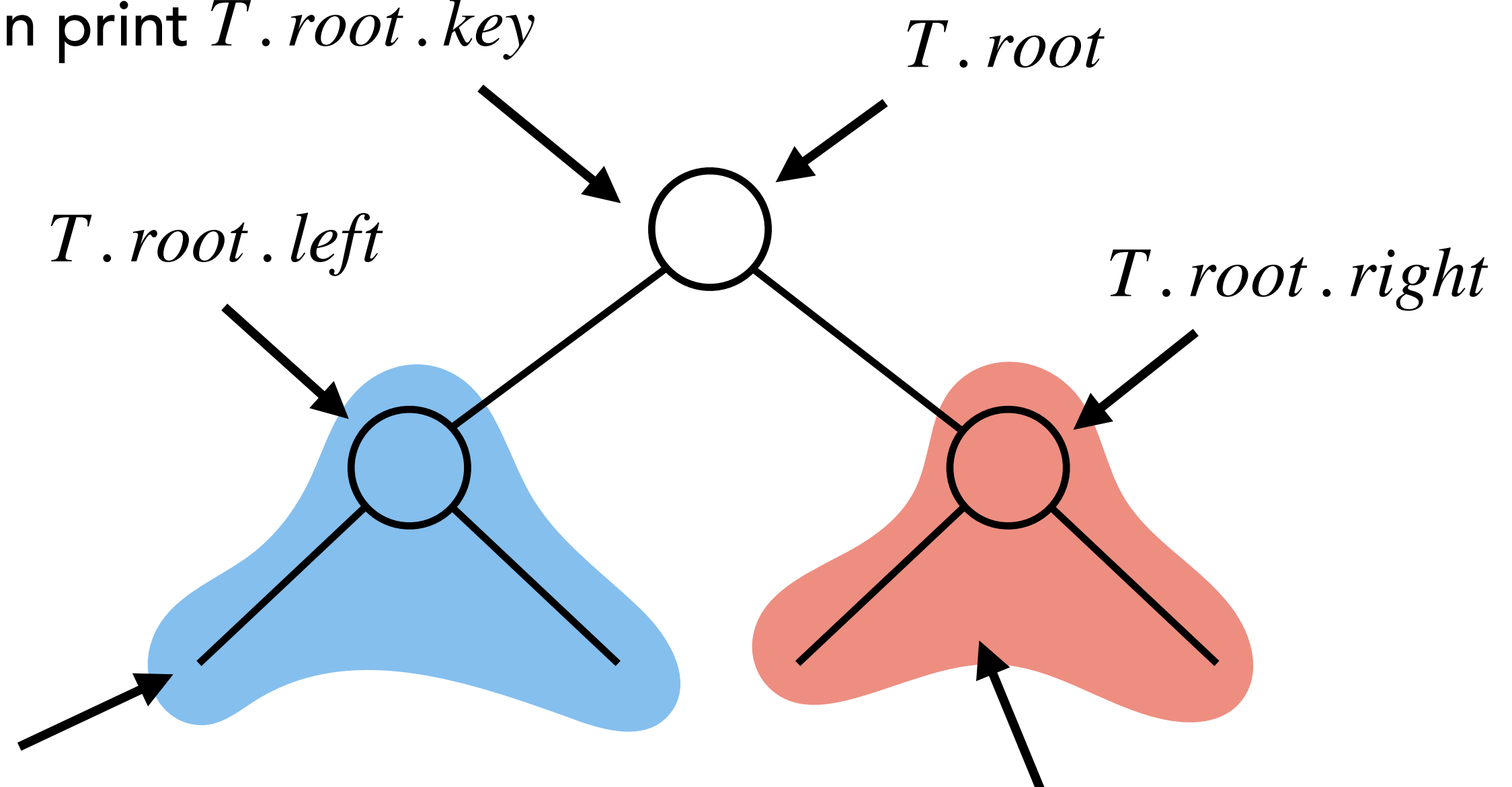
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Inorder Tree Walk: Time Analysis

Inorder-Tree-Walk(x):

1. **if $x \neq \text{NIL}$**
2. **Inorder-Tree-Walk($x . \text{left}$)**
3. **print $x . \text{key}$**
4. **Inorder-Tree-Walk($x . \text{right}$)**

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Runtime: $\Theta(n)$, where n is the # of nodes in the tree.

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Can be proven using induction.



Search in a BST

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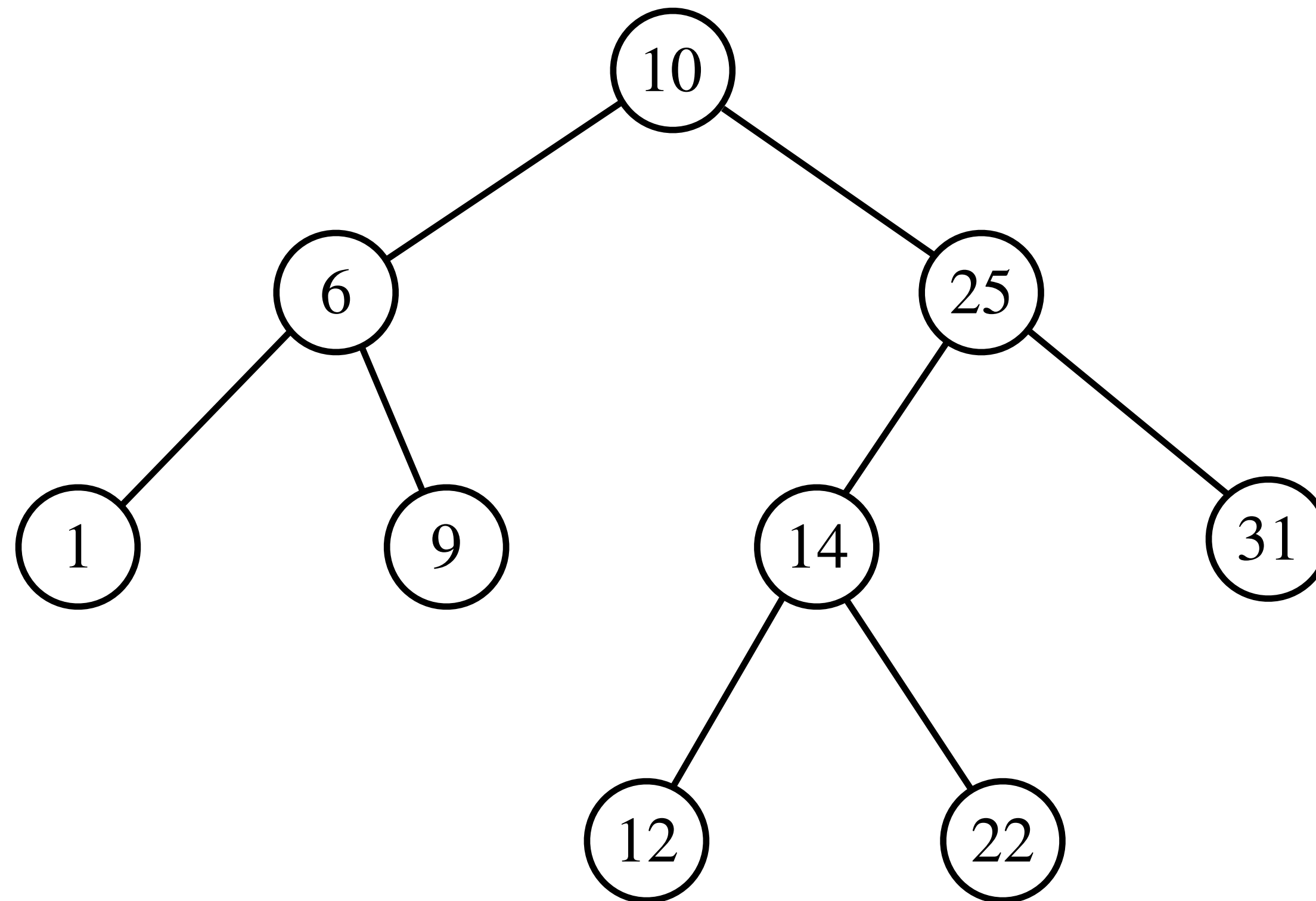
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Illustration: Searching for 12 in the below BST.

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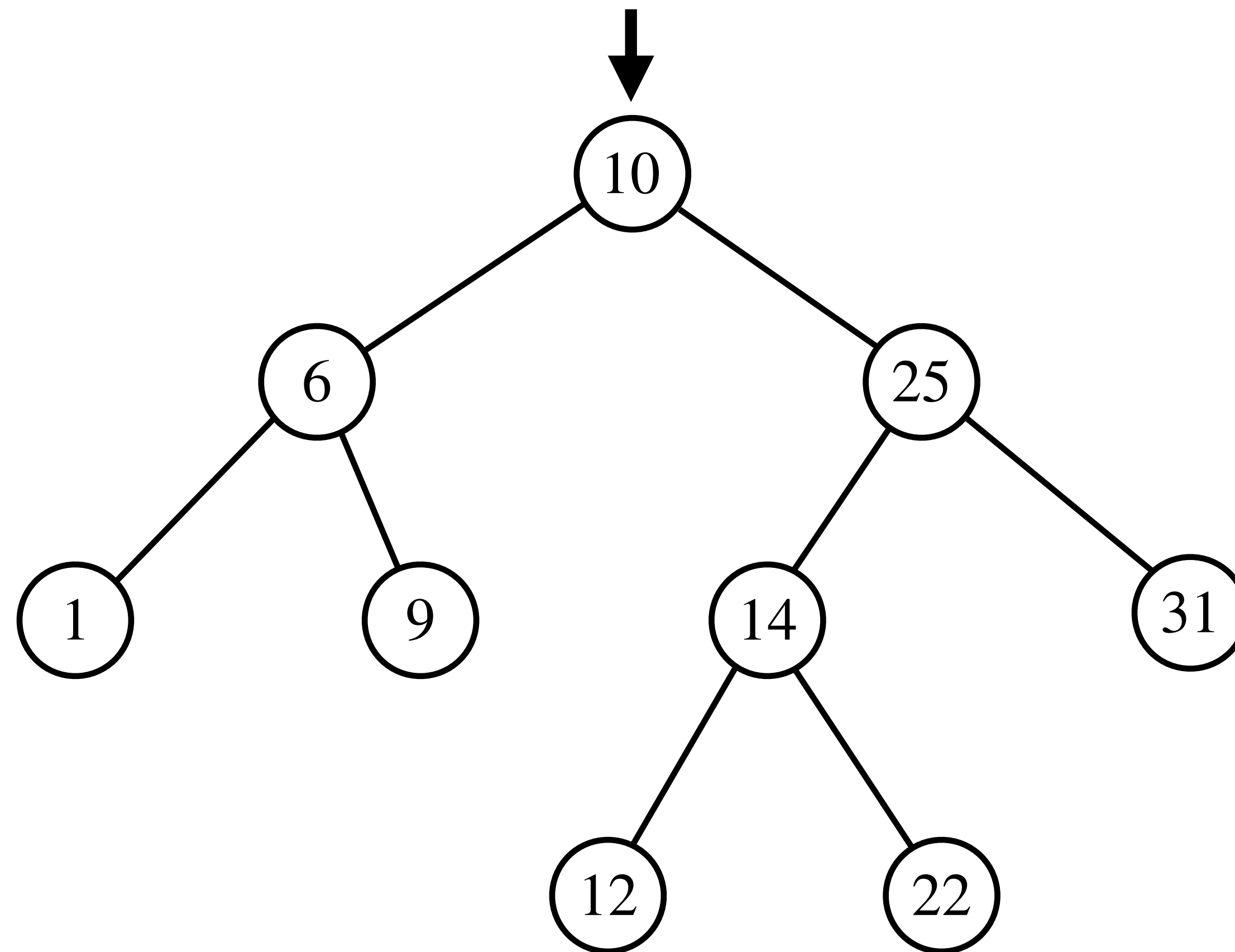
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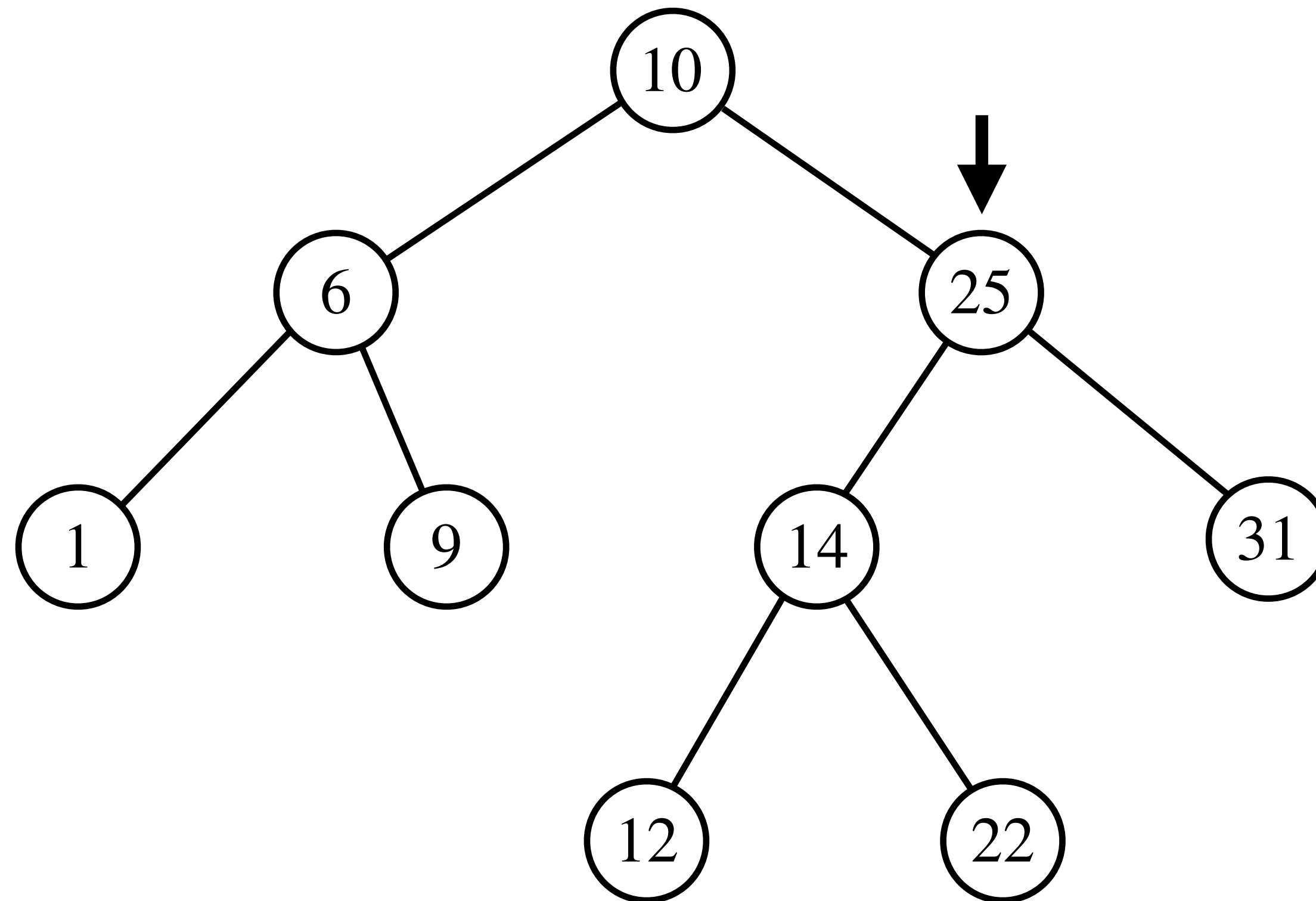
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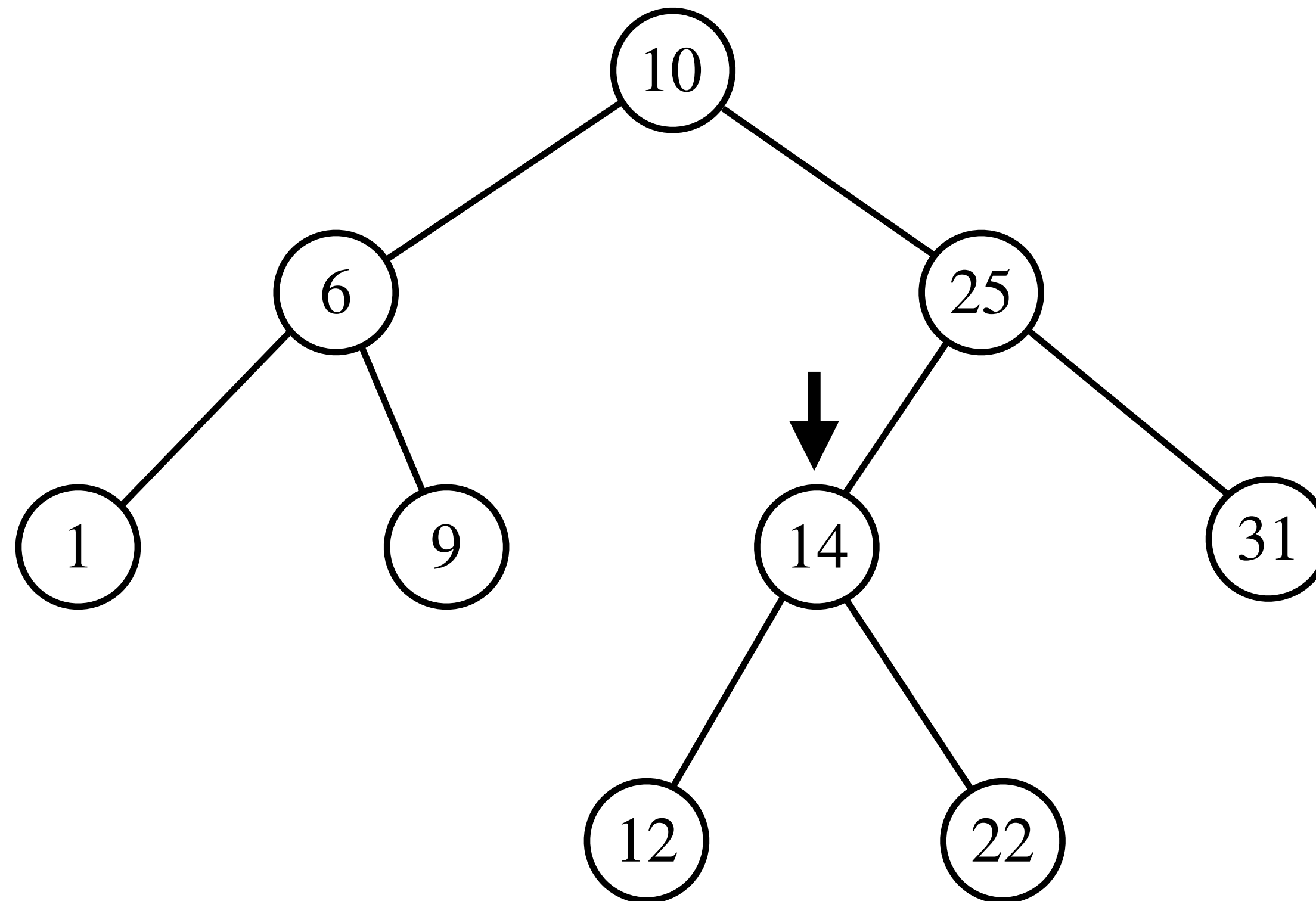
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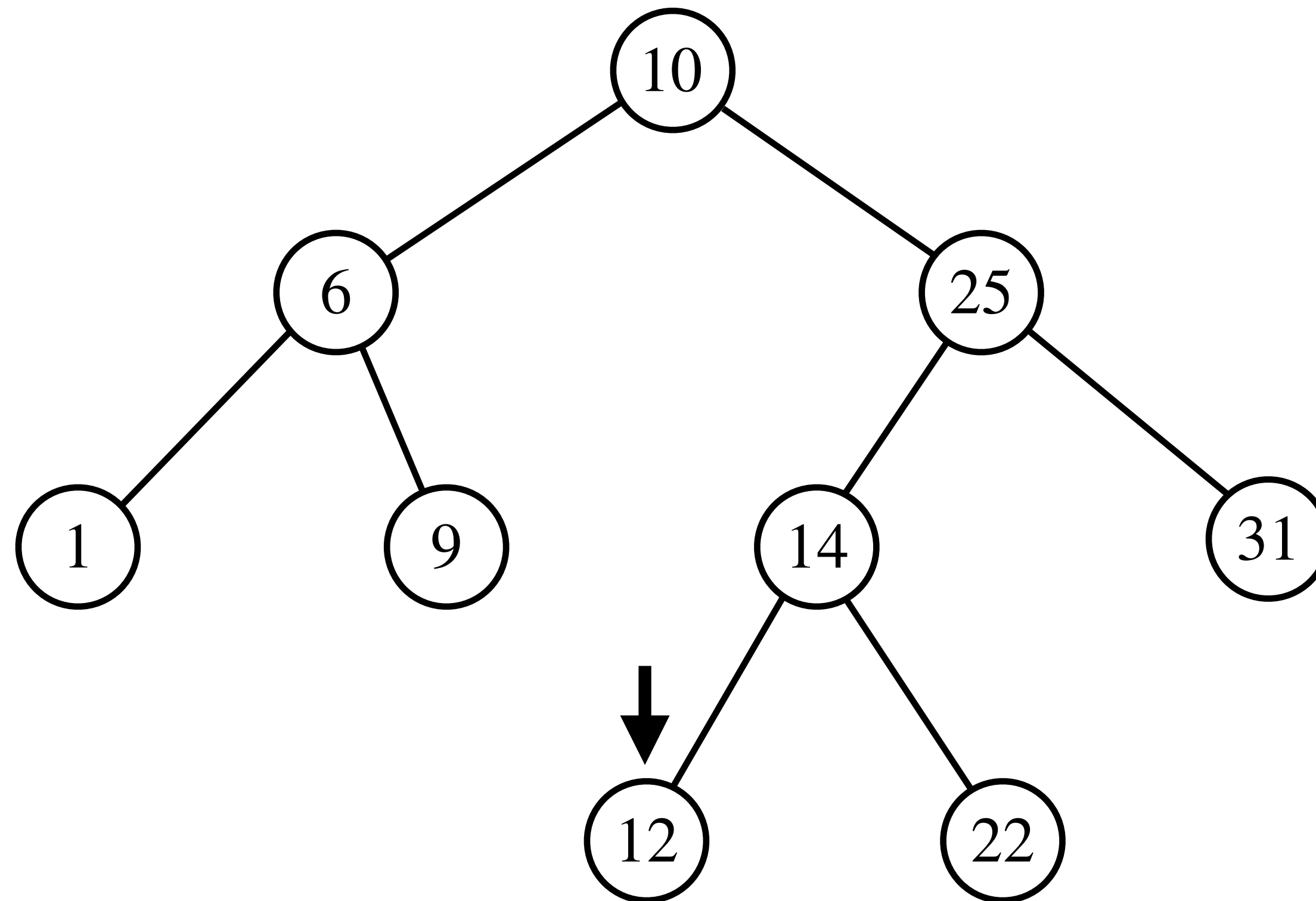
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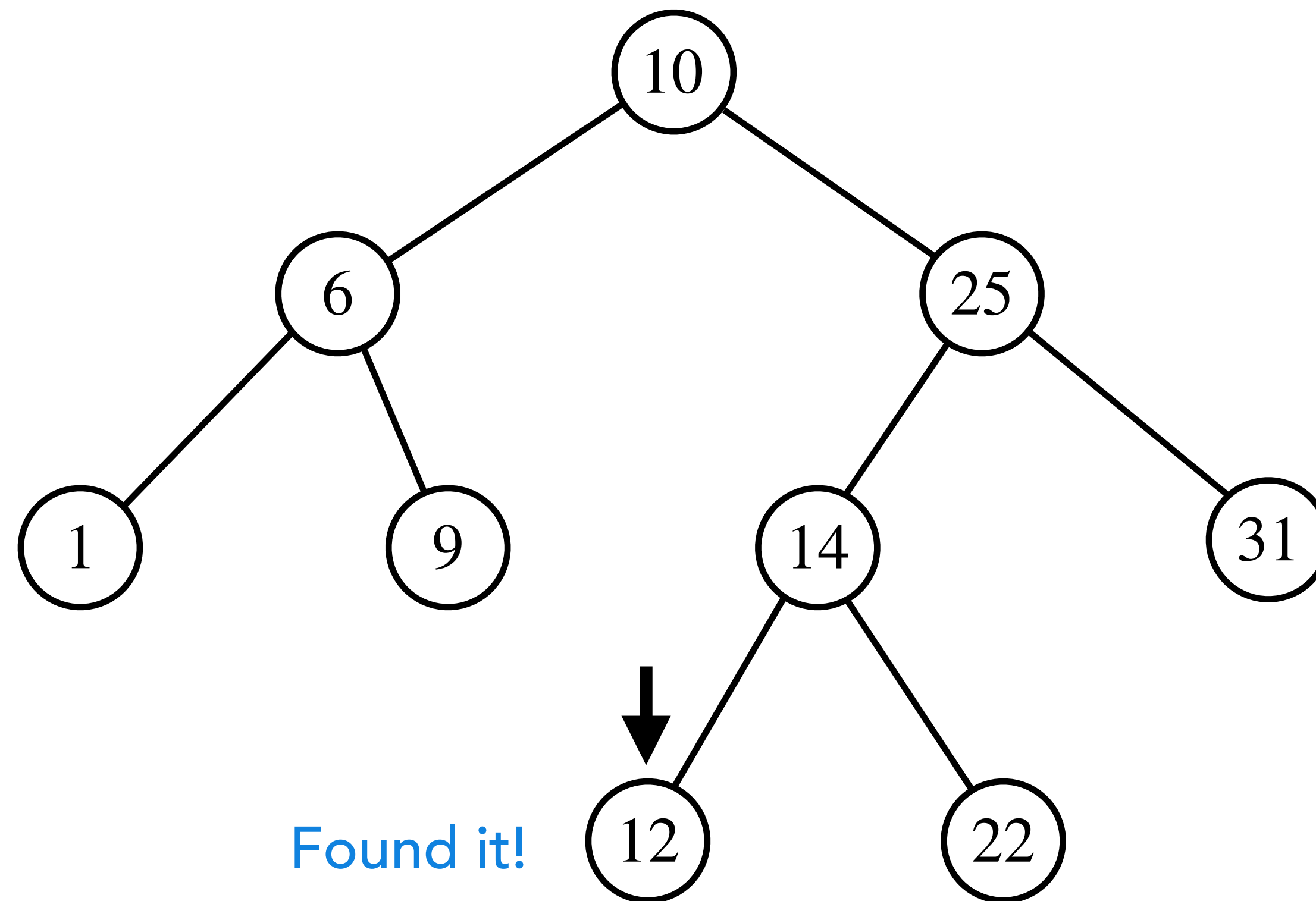
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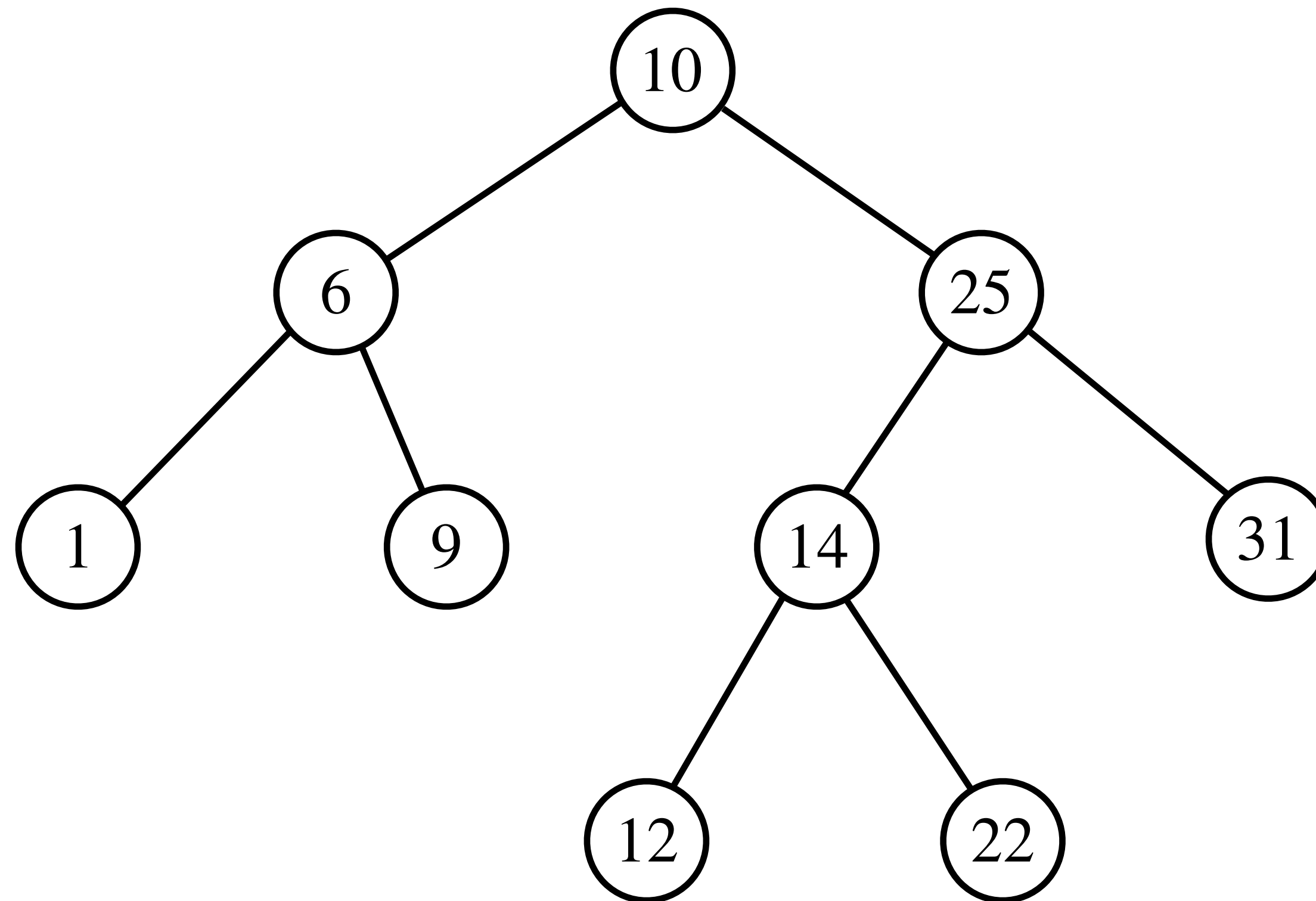
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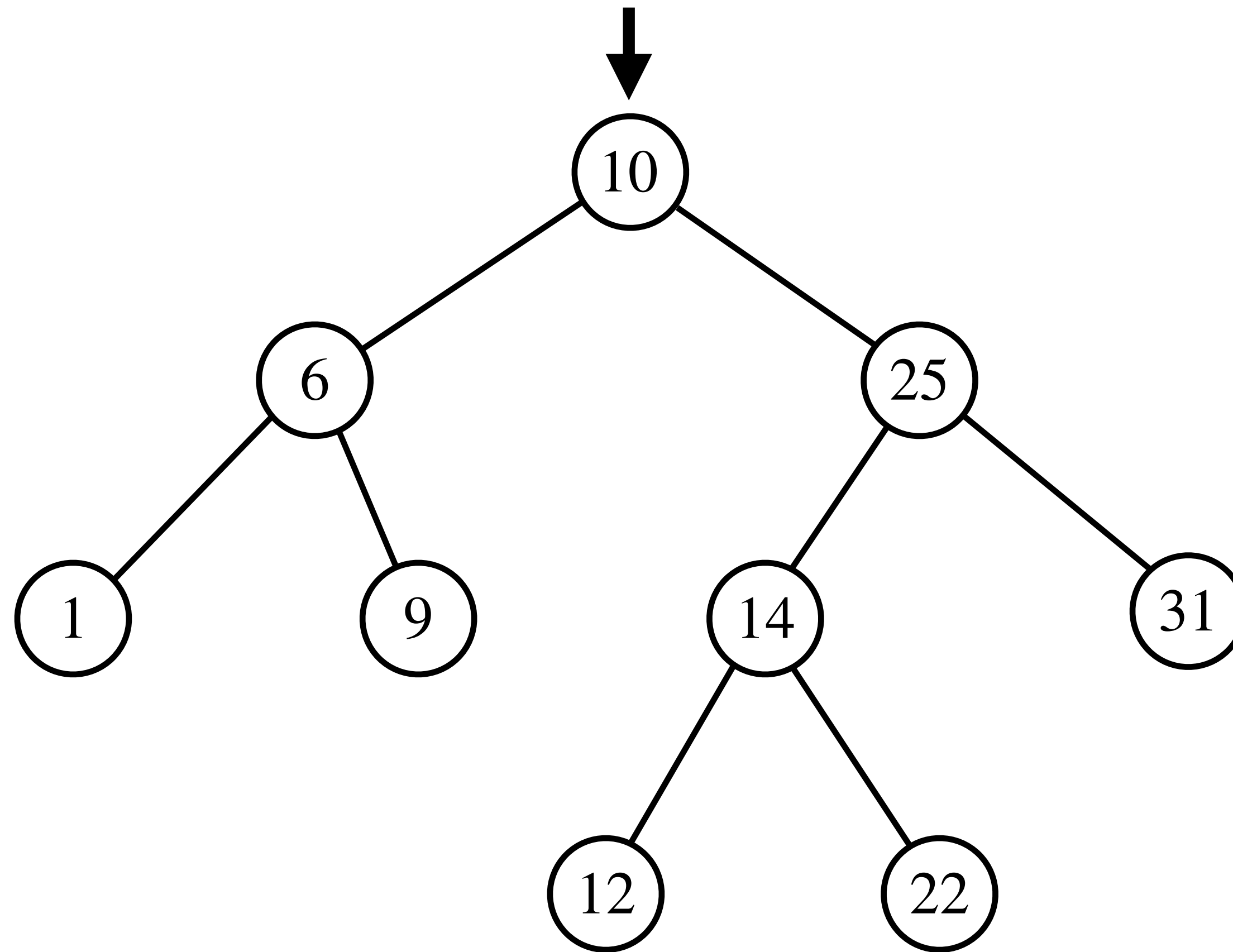
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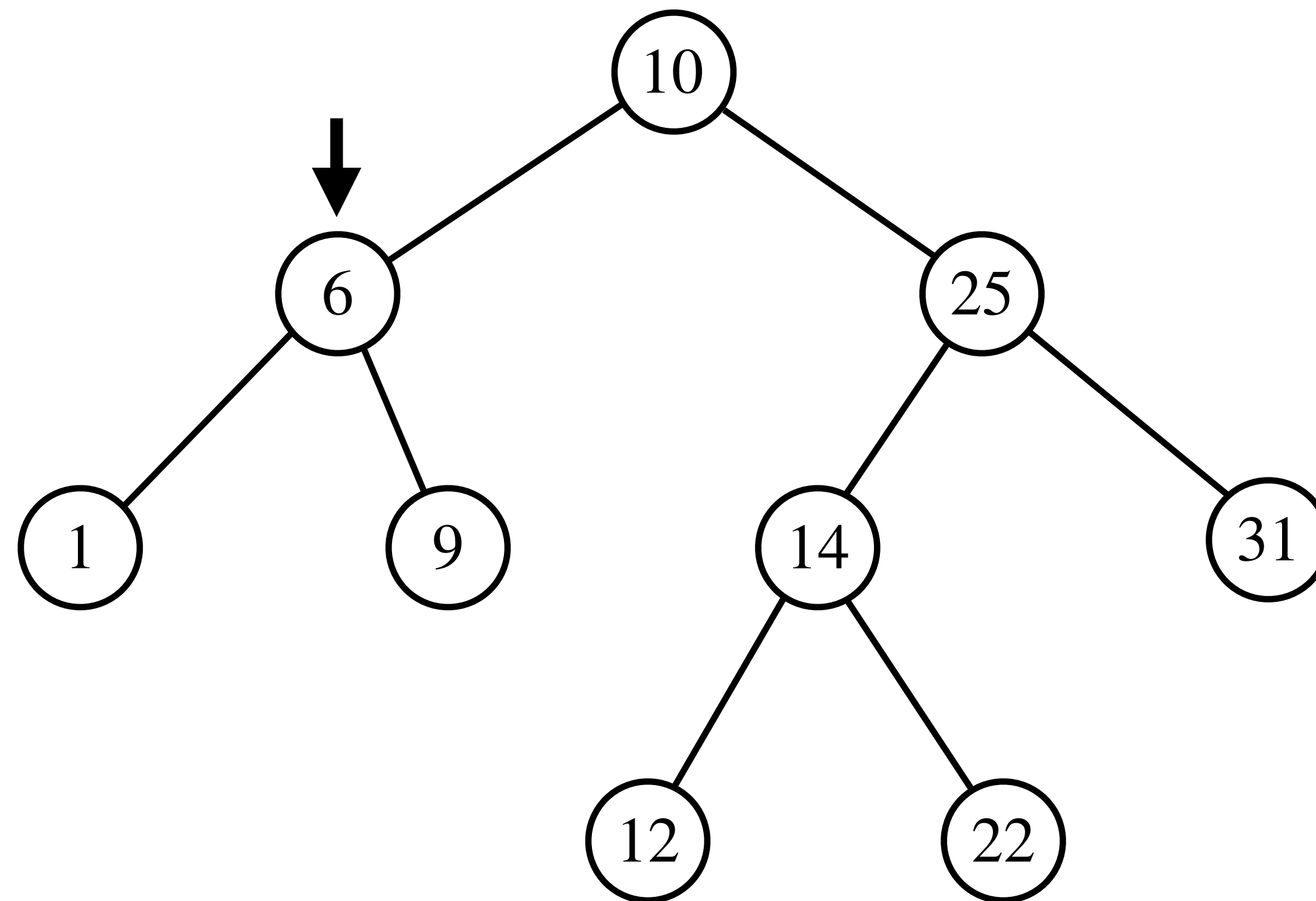
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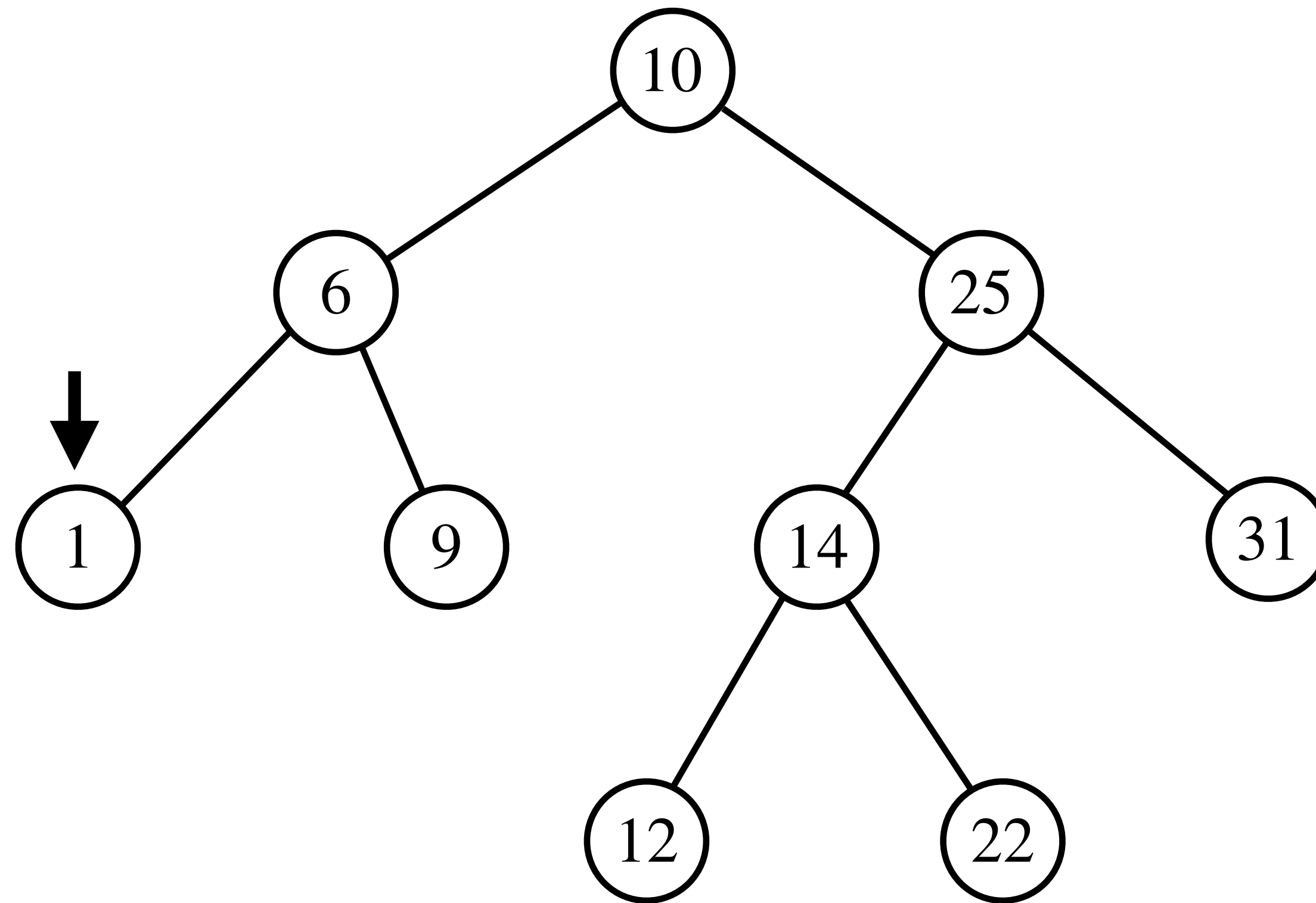
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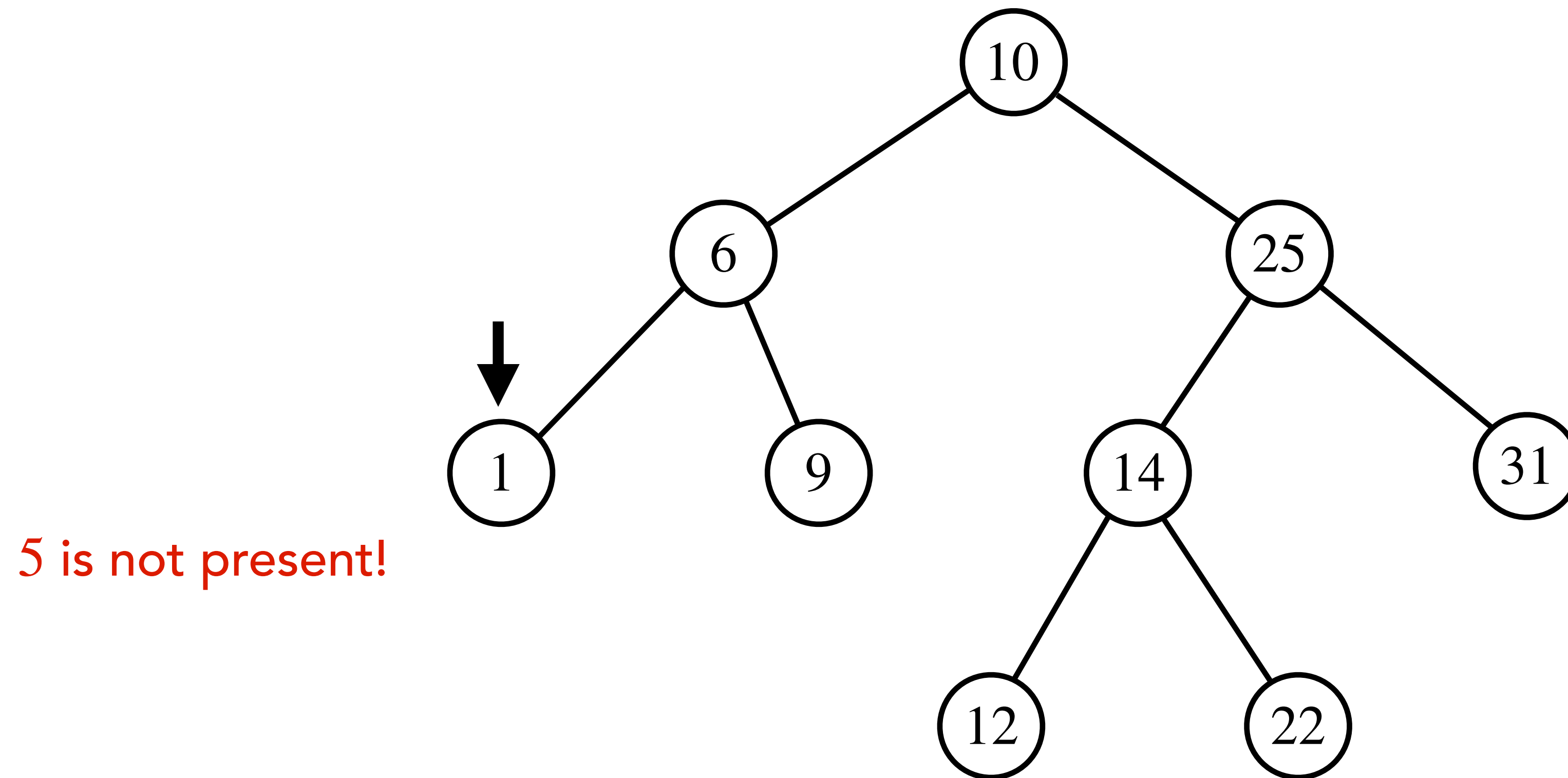
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$$k = 8$$

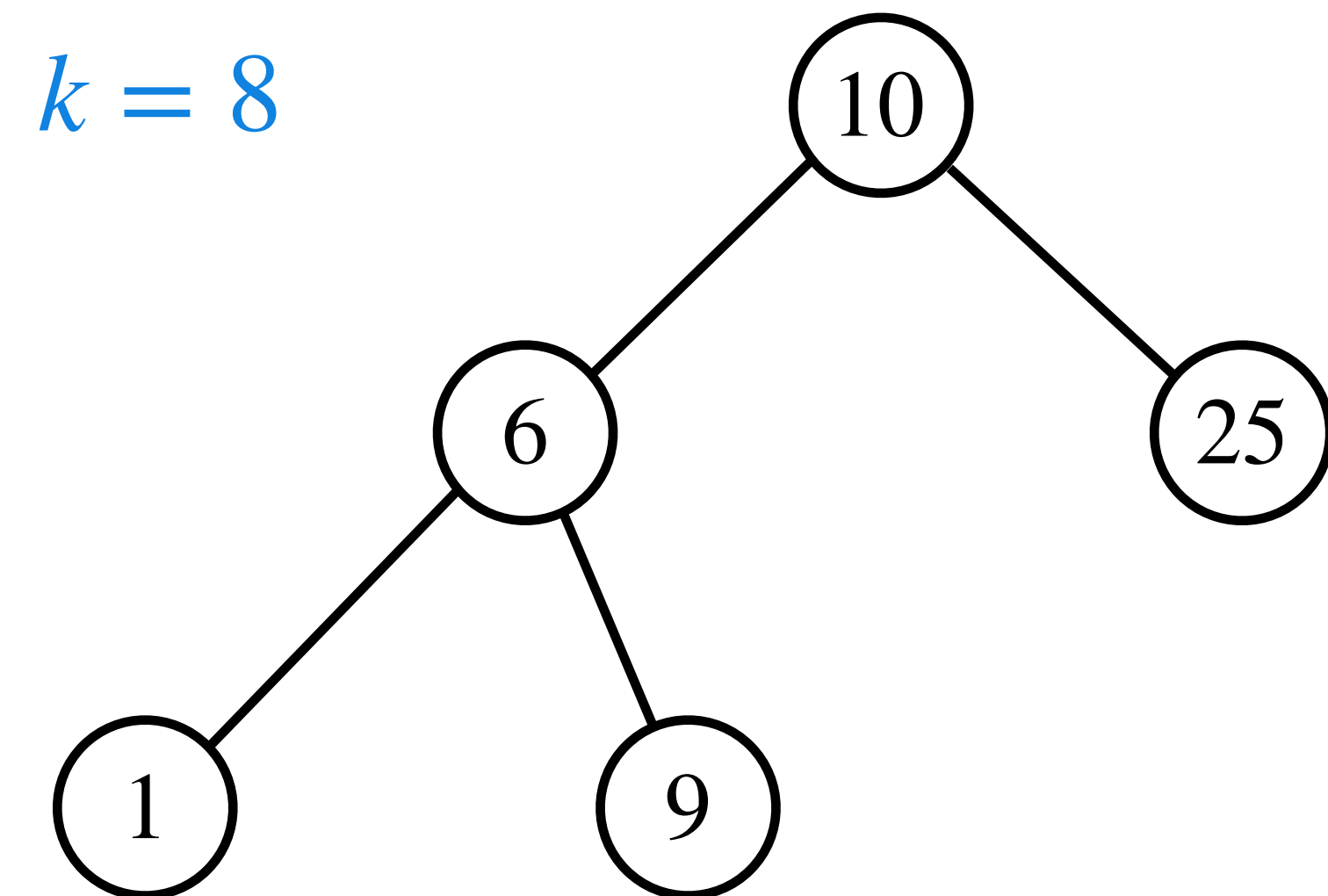
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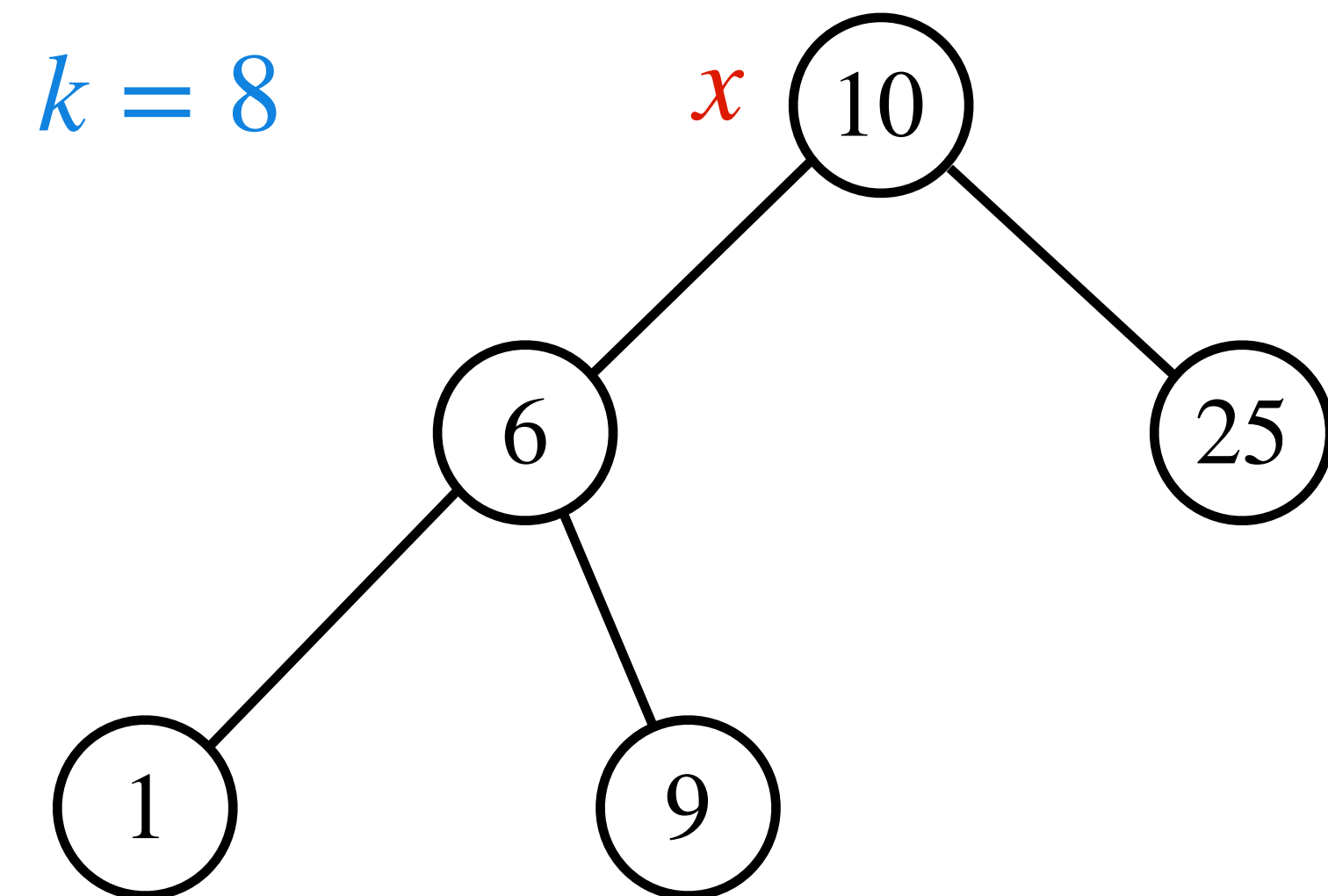
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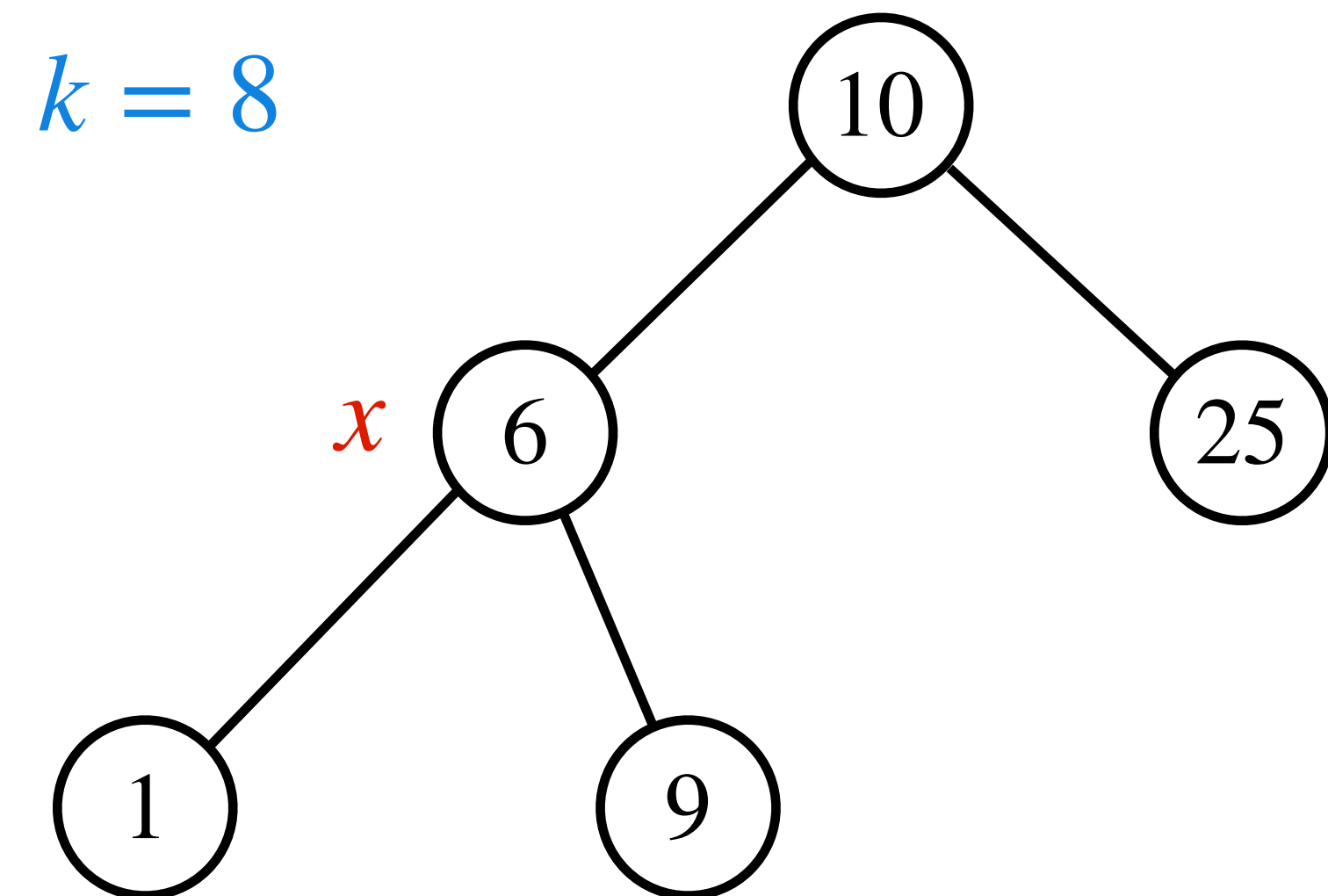
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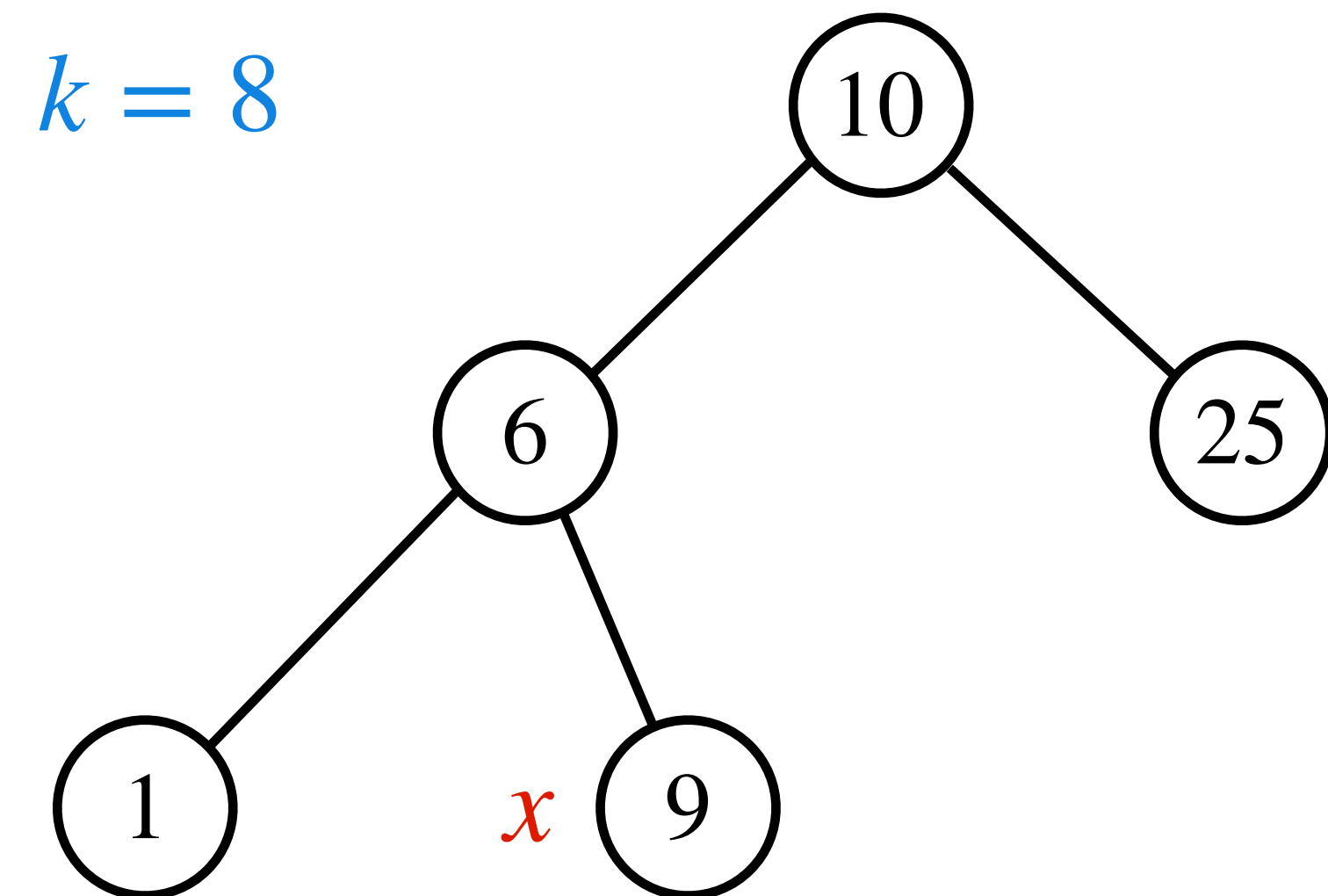
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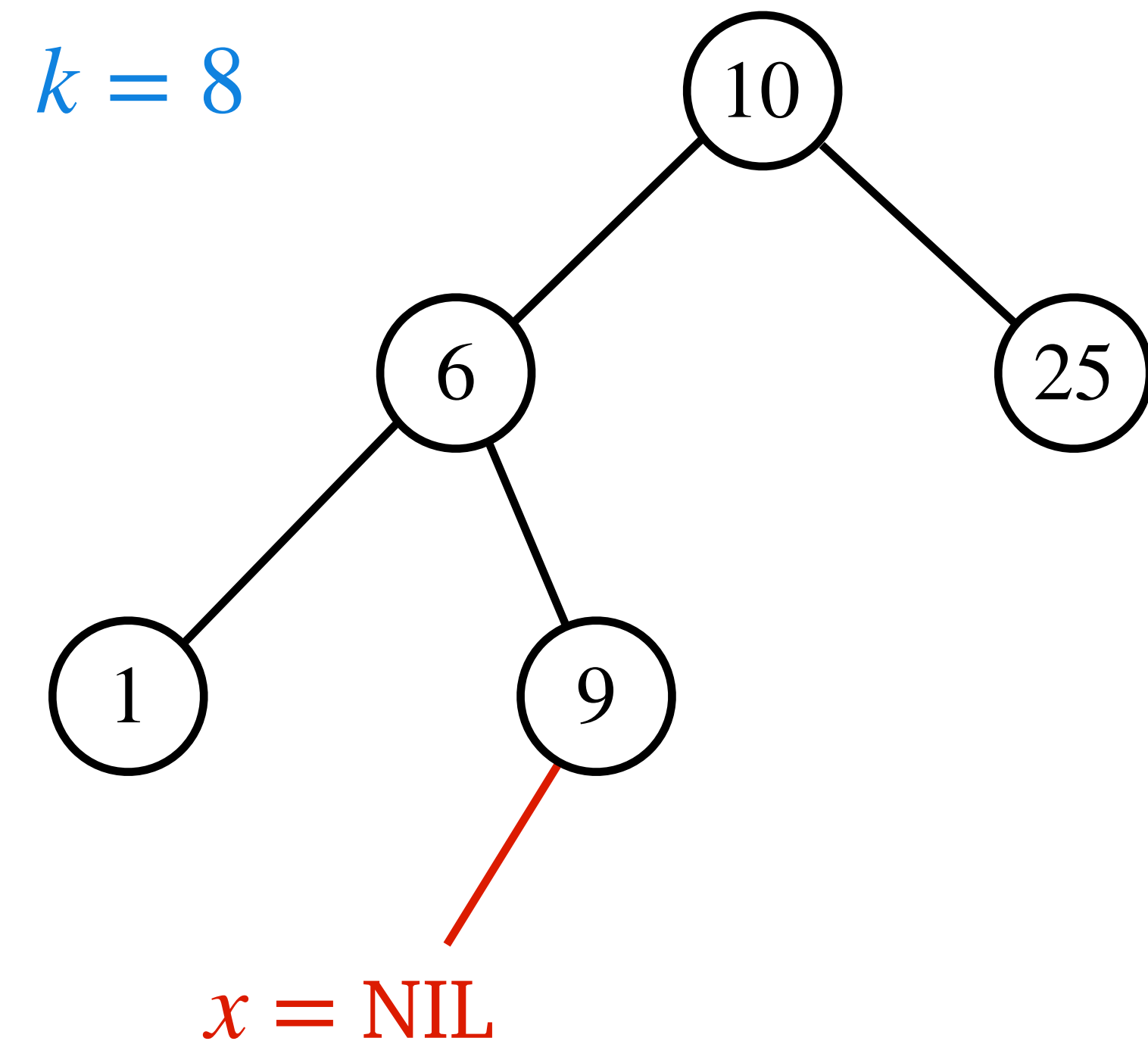
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Runtime: $\Theta(h)$, where h = height of T , as while loop goes one level down with every iteration.

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Illustration:

Finding Successor in a BST

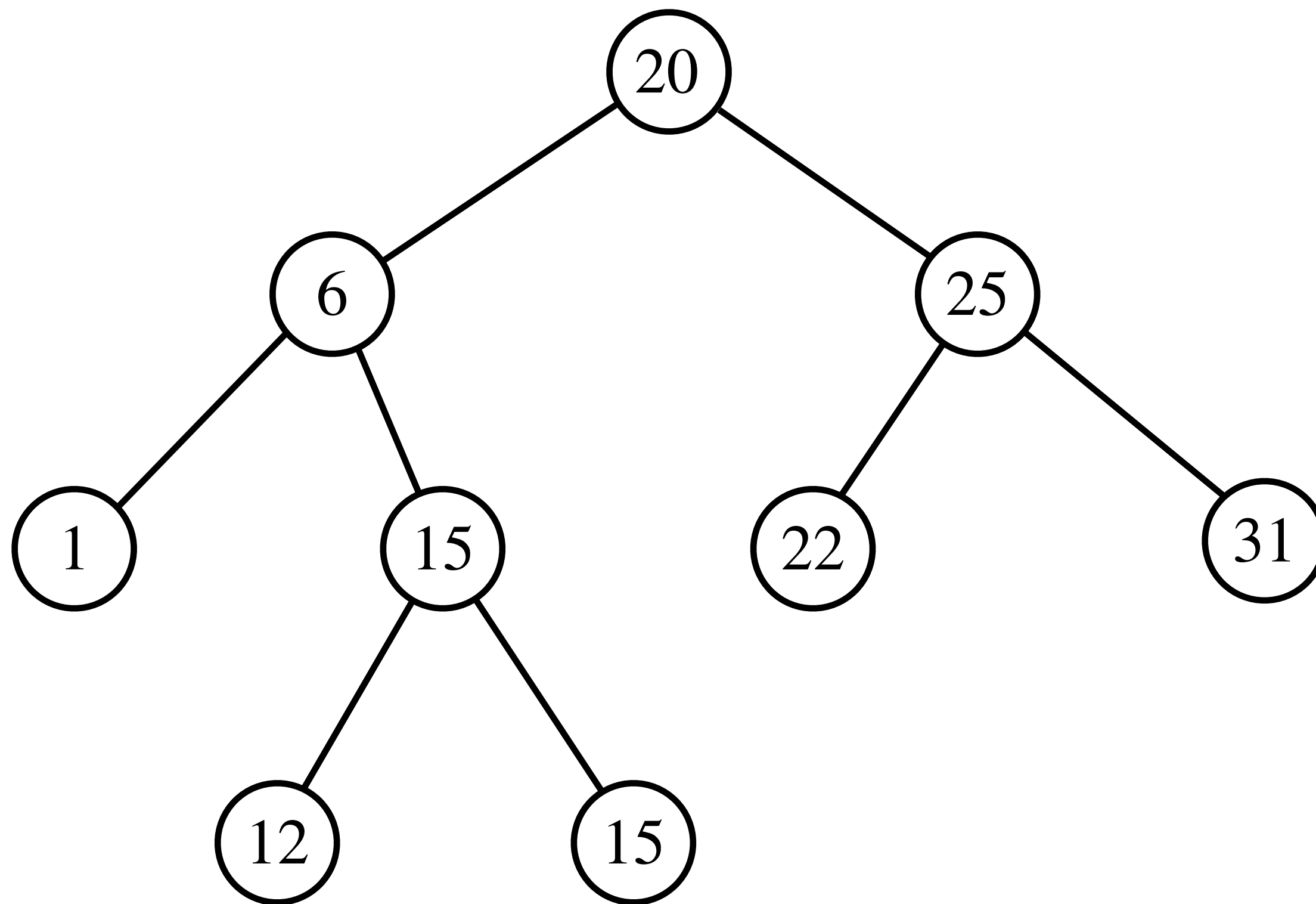
Goal: Given a **node** x of a BST find its **successor**.  Node printed after x in inorder-walk.

Illustration: Find the successor of 6 in the below BST.

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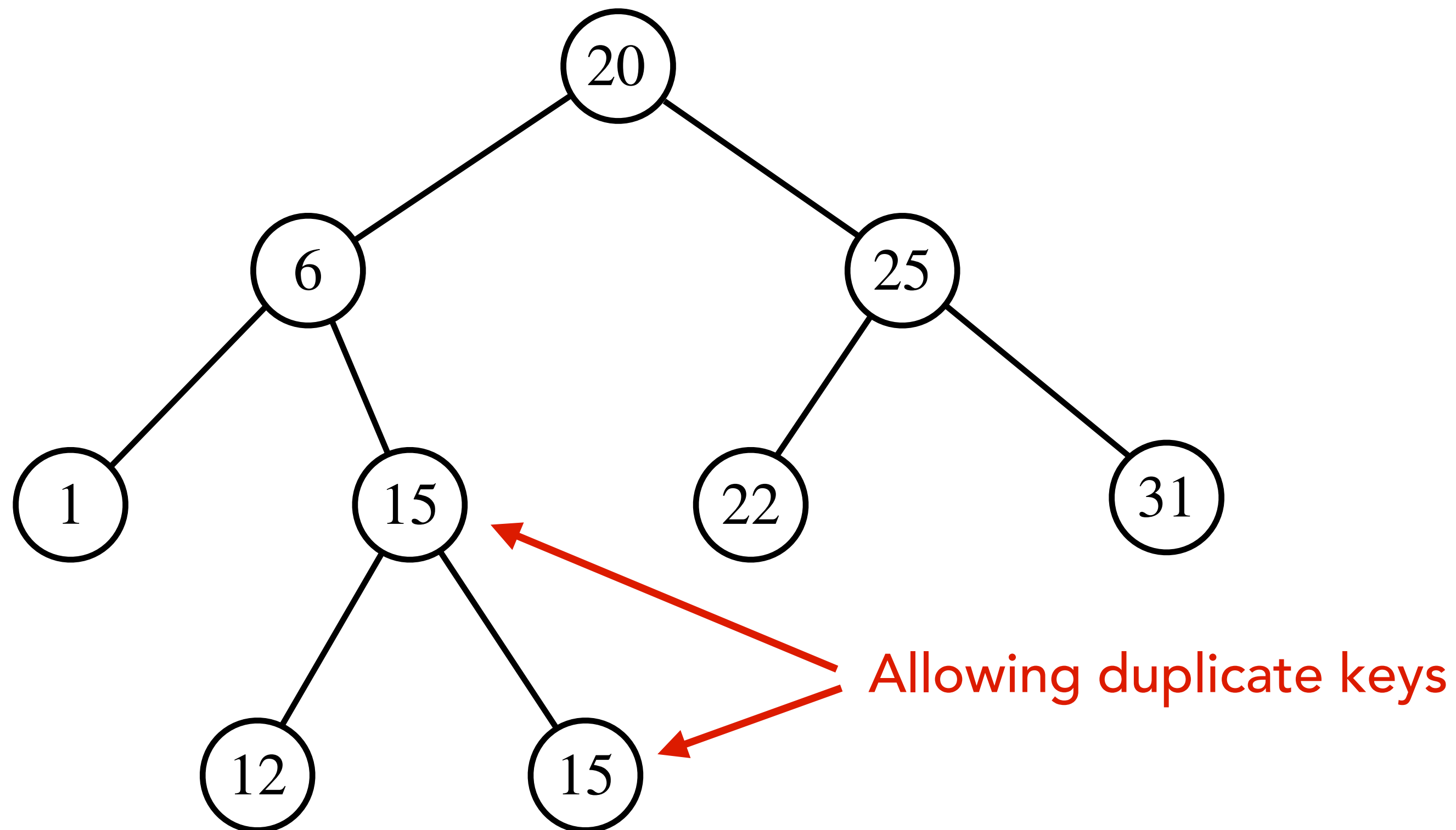
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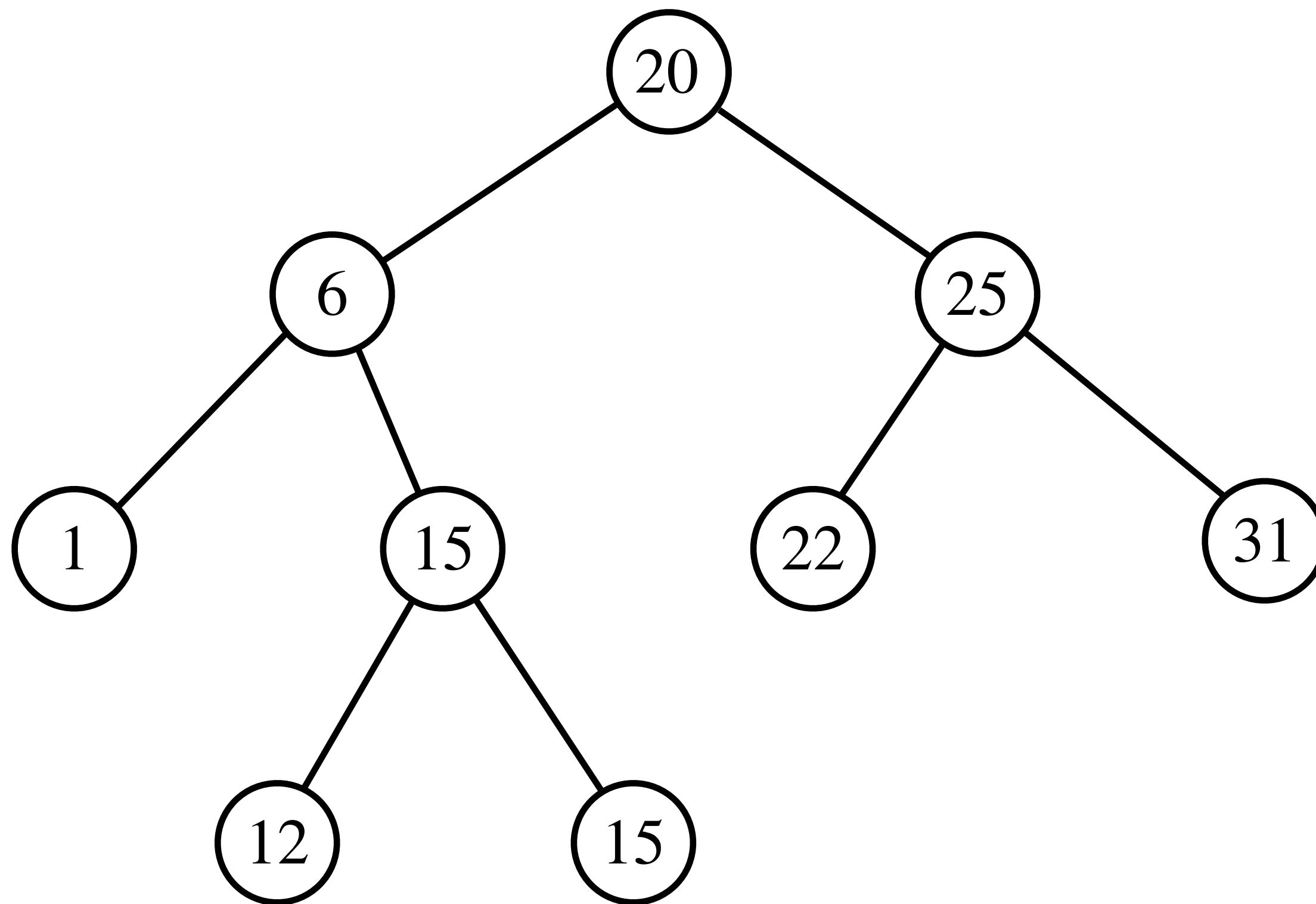
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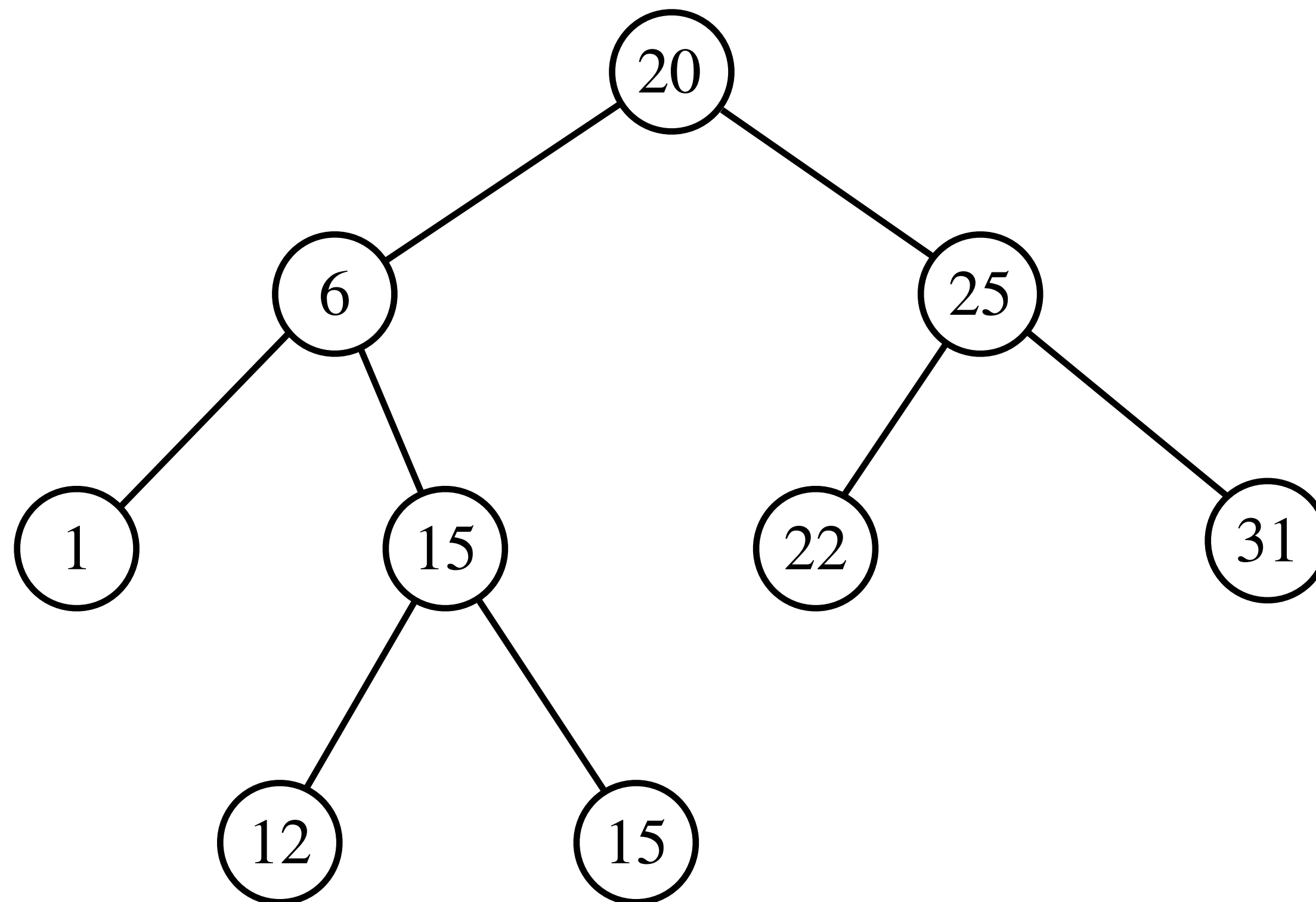
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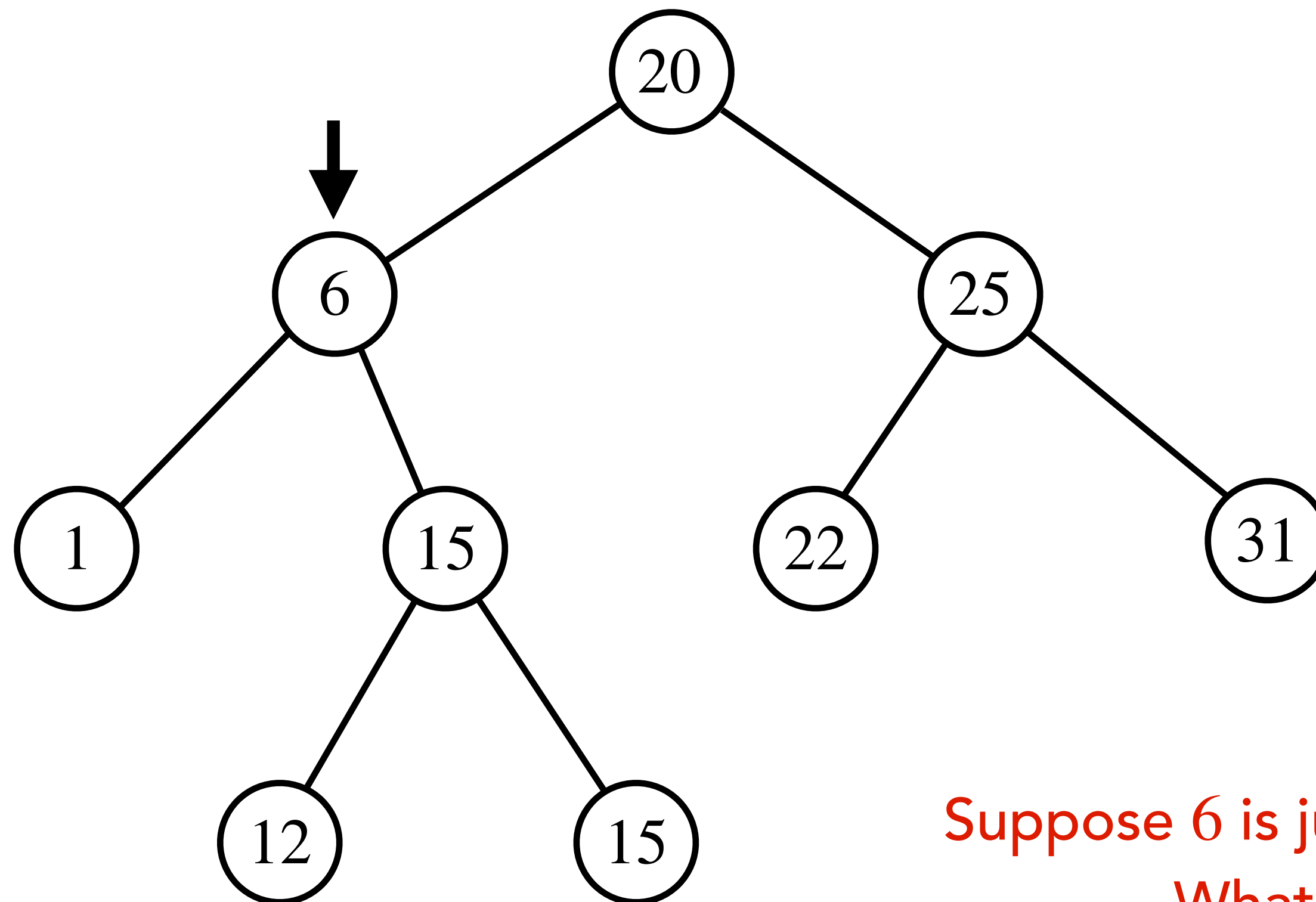
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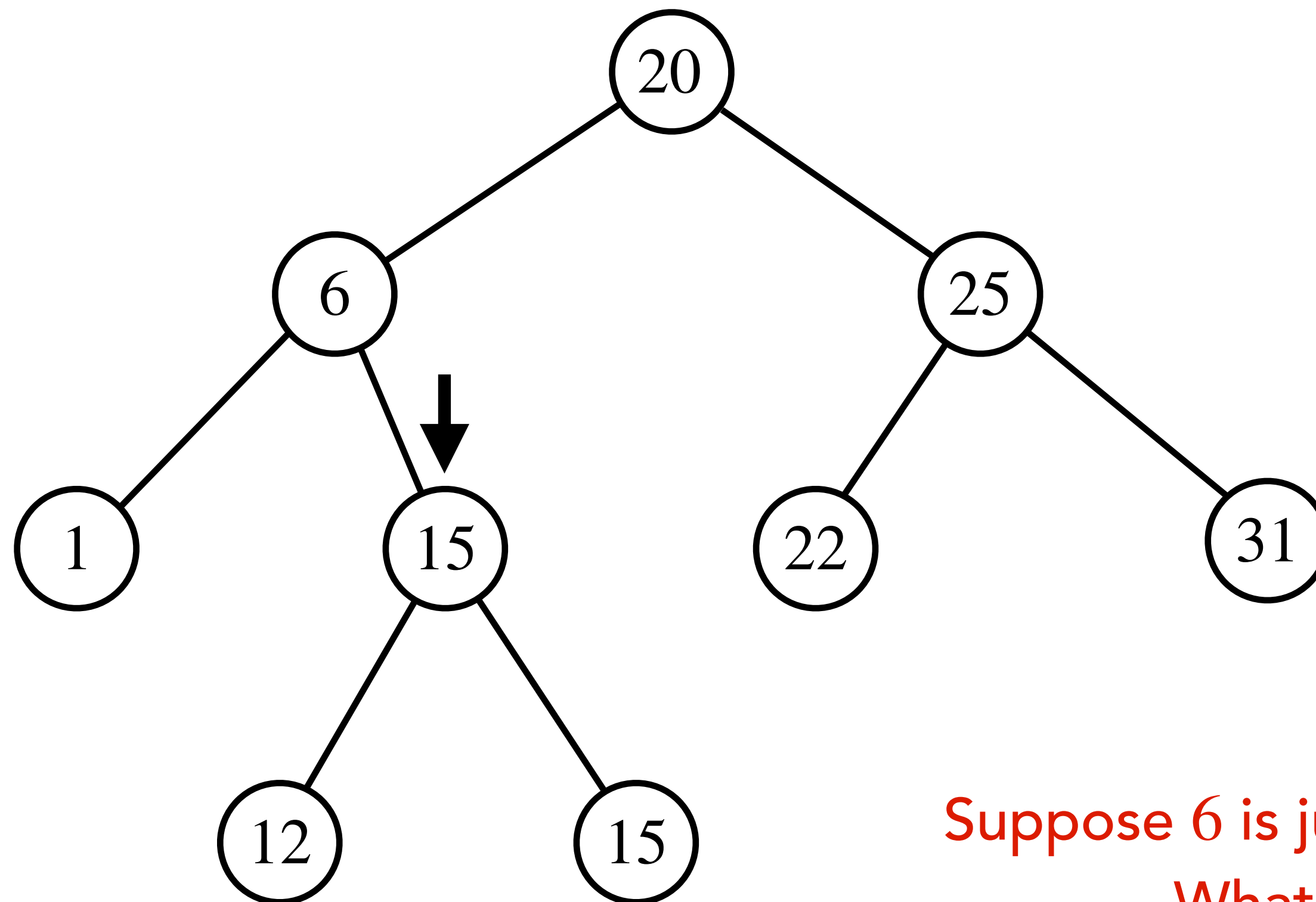
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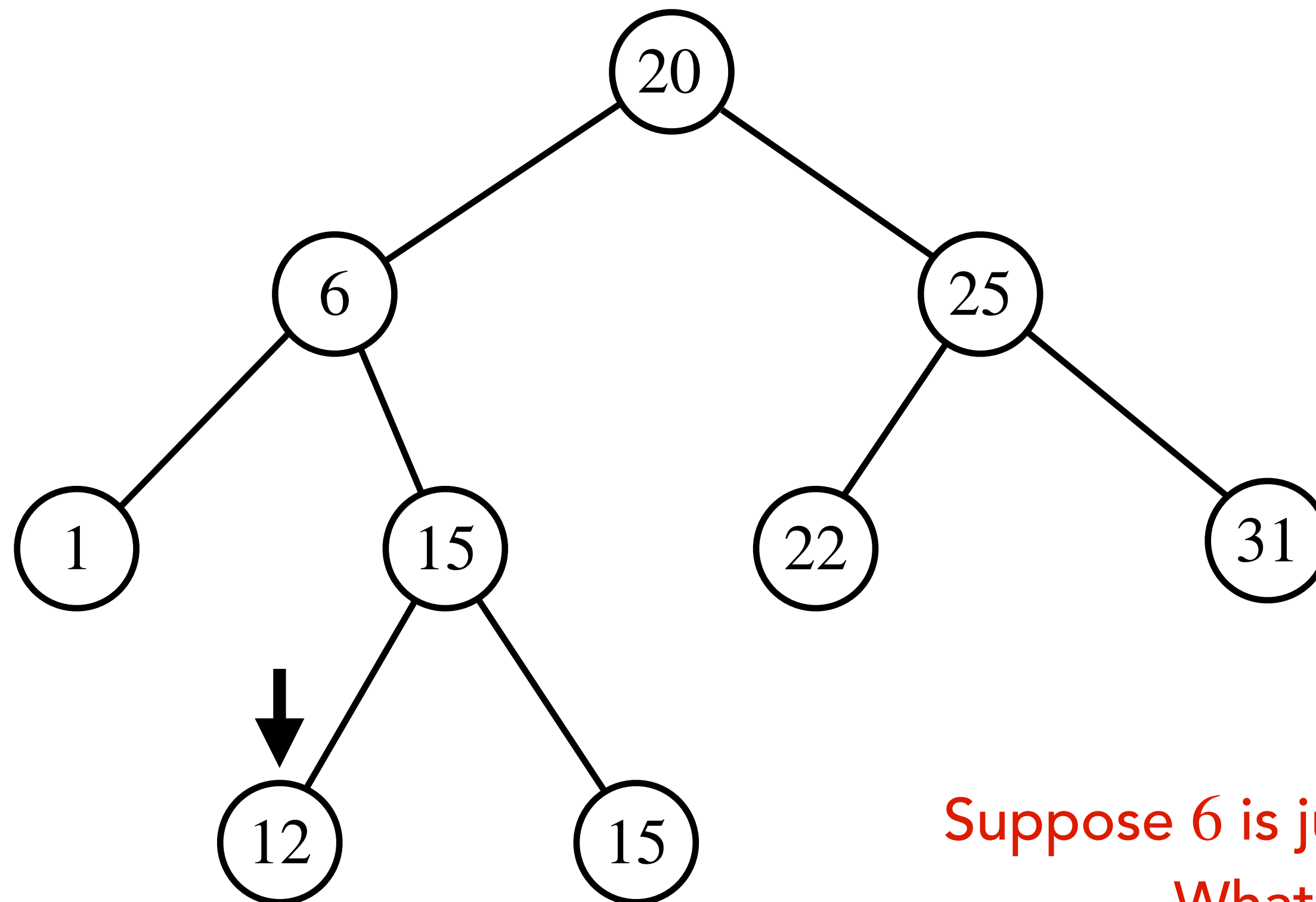
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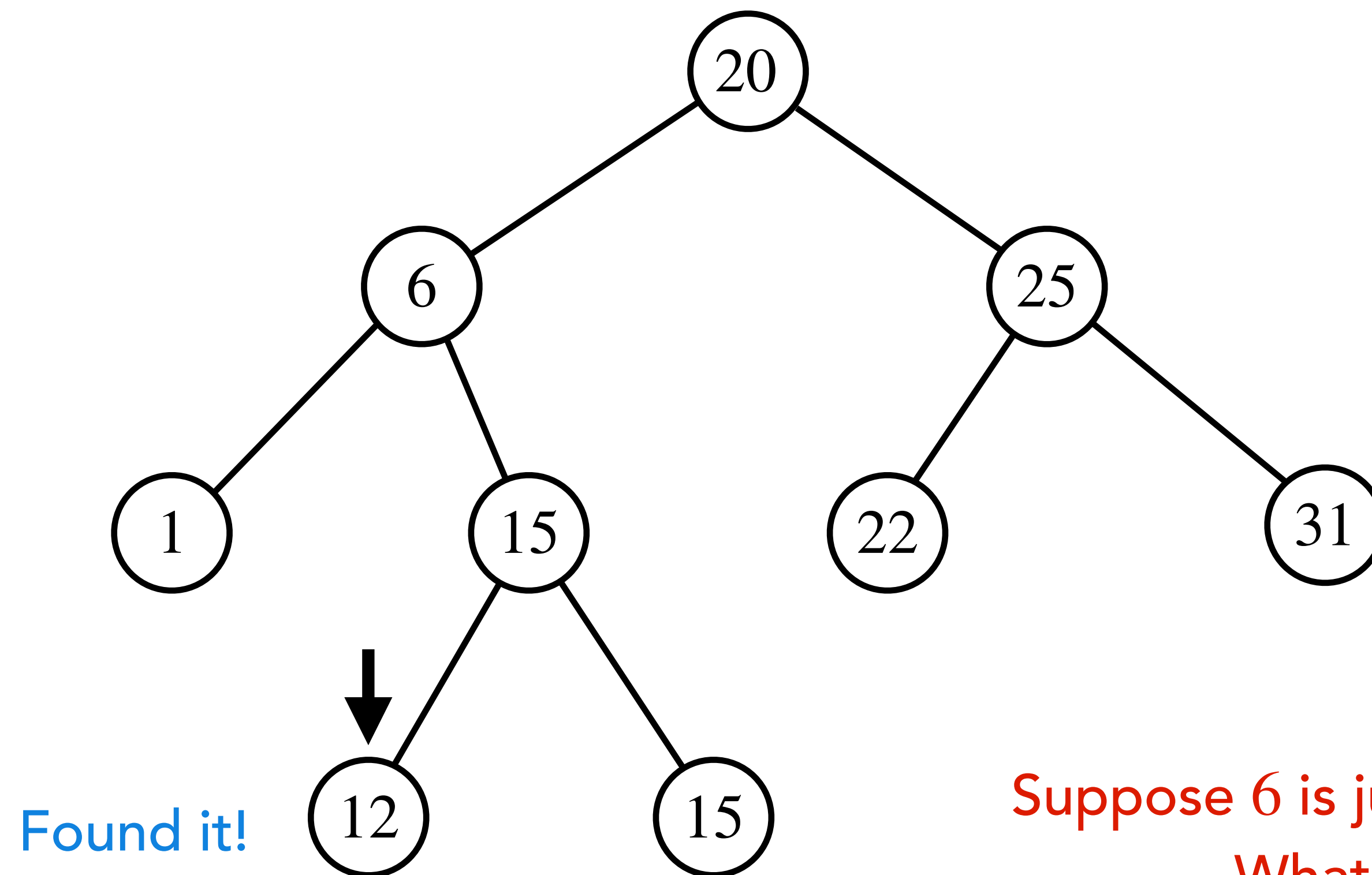
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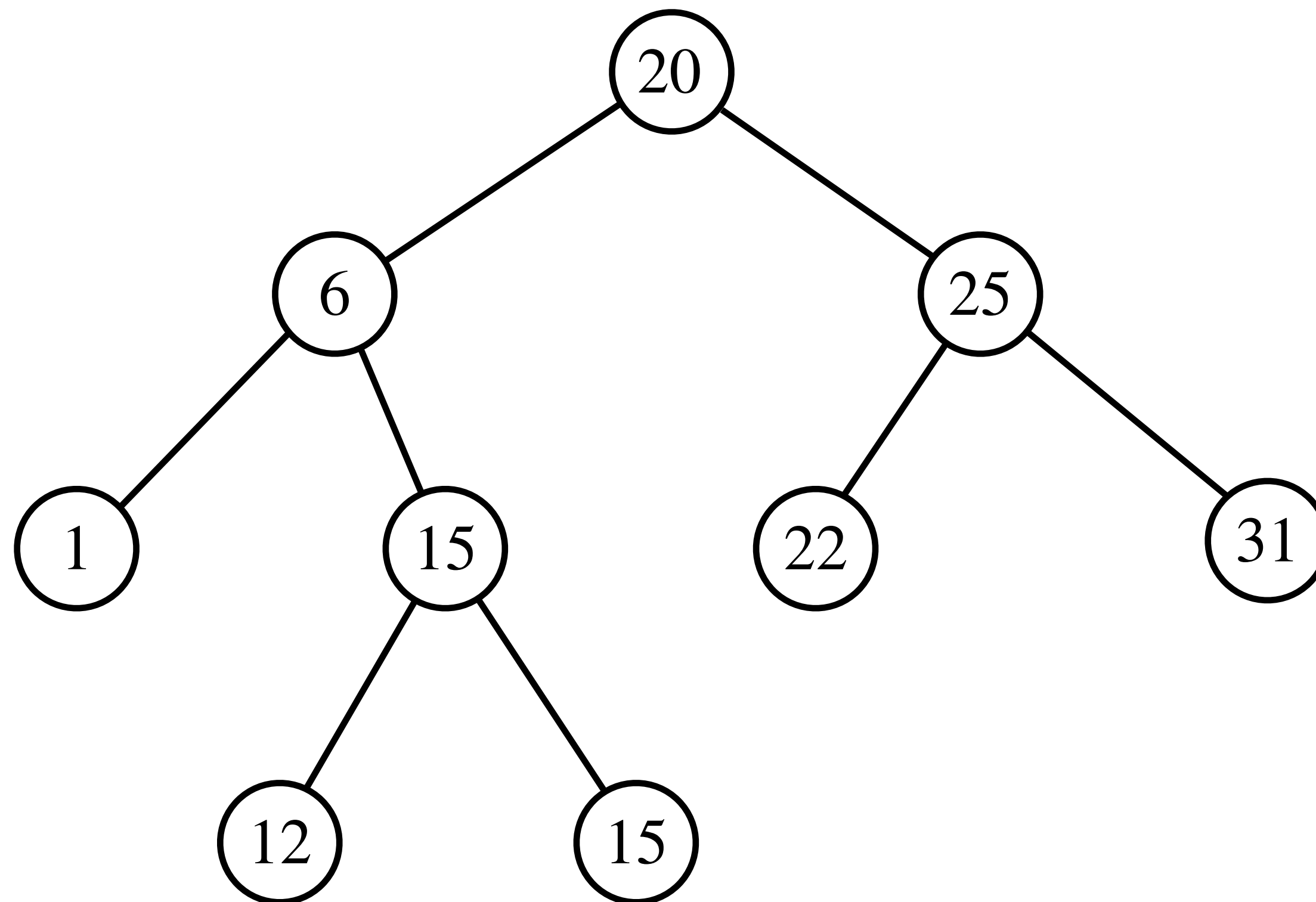
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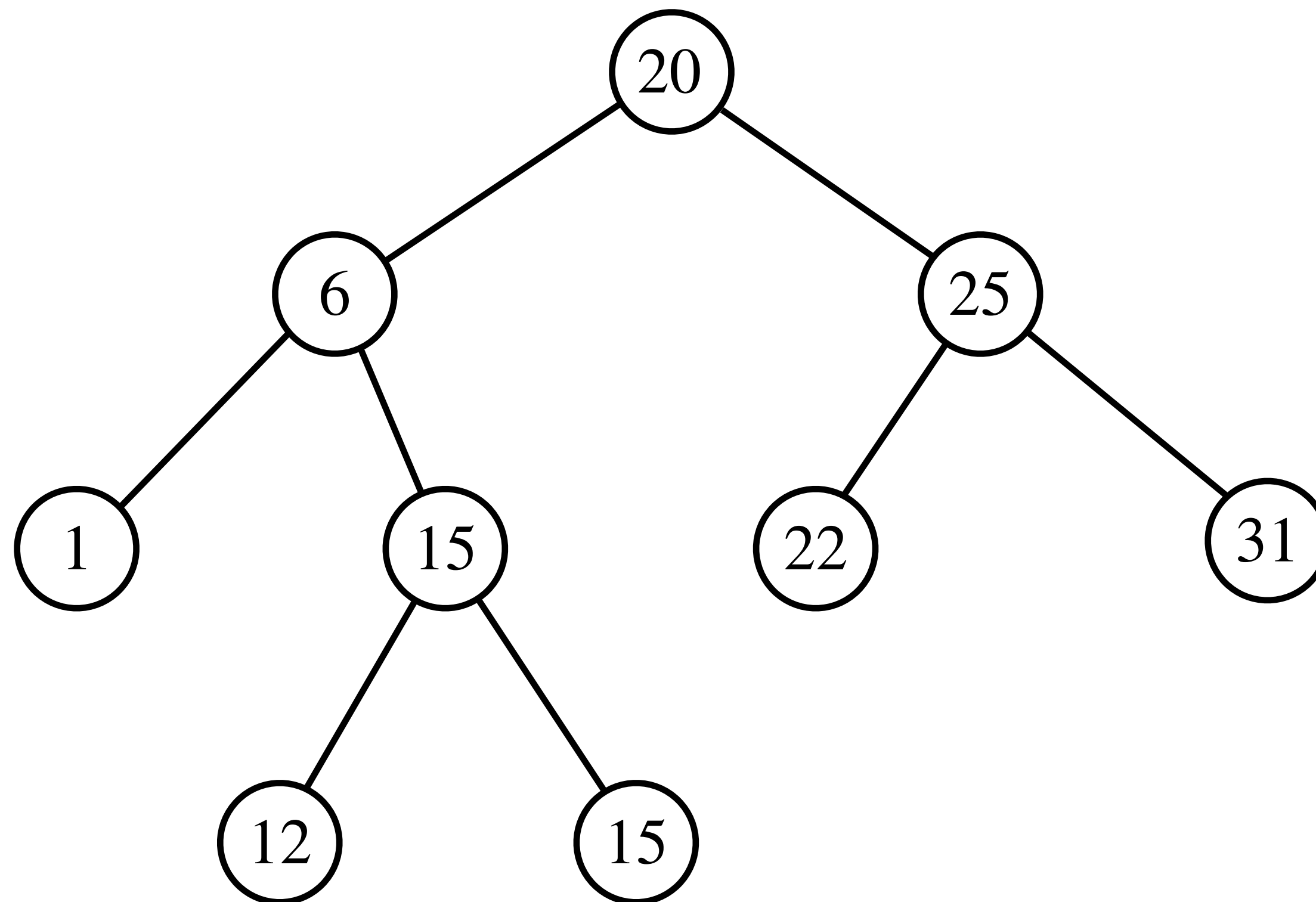
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Illustration: Find the successor of the leaf 15 in the below BST.



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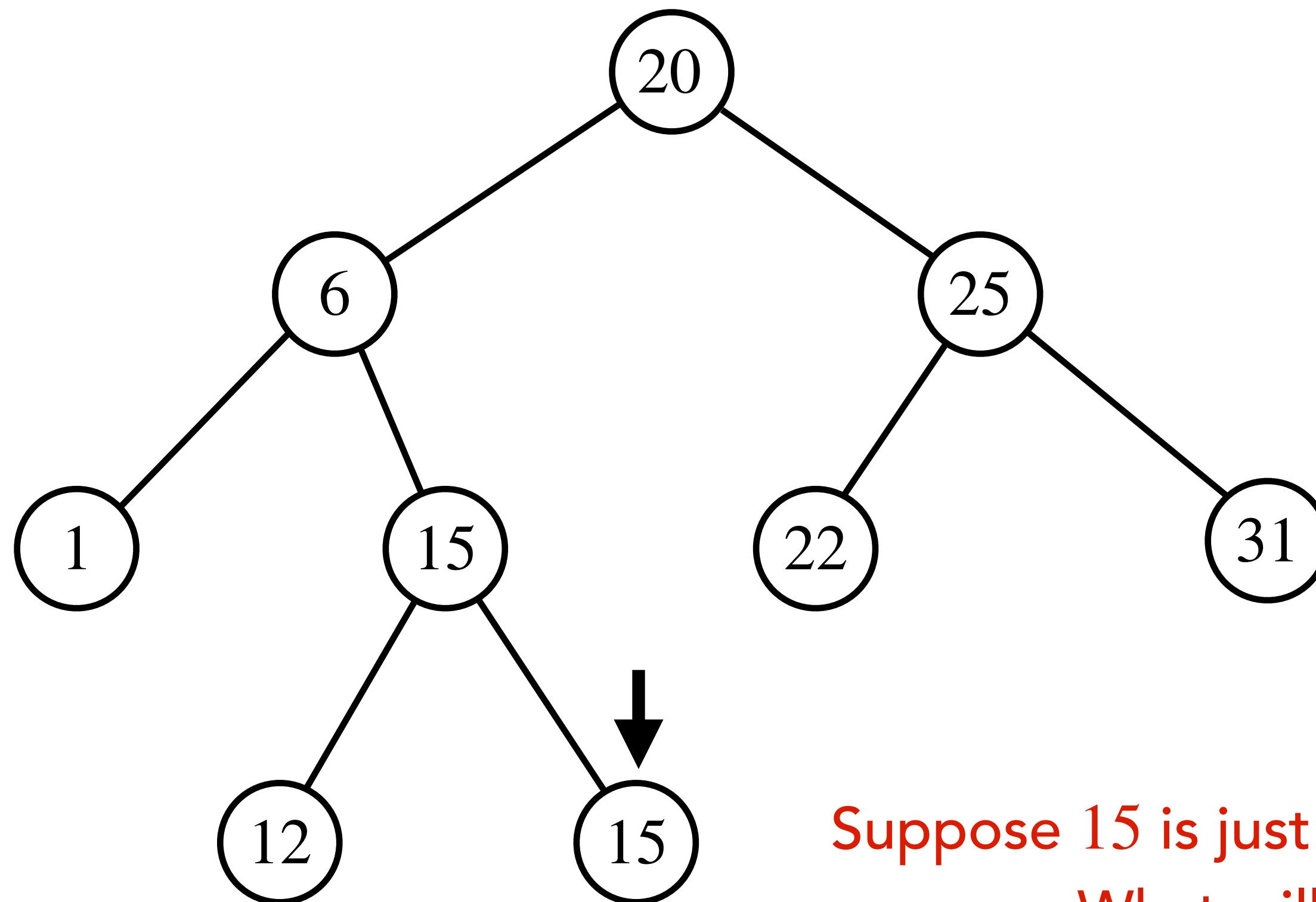
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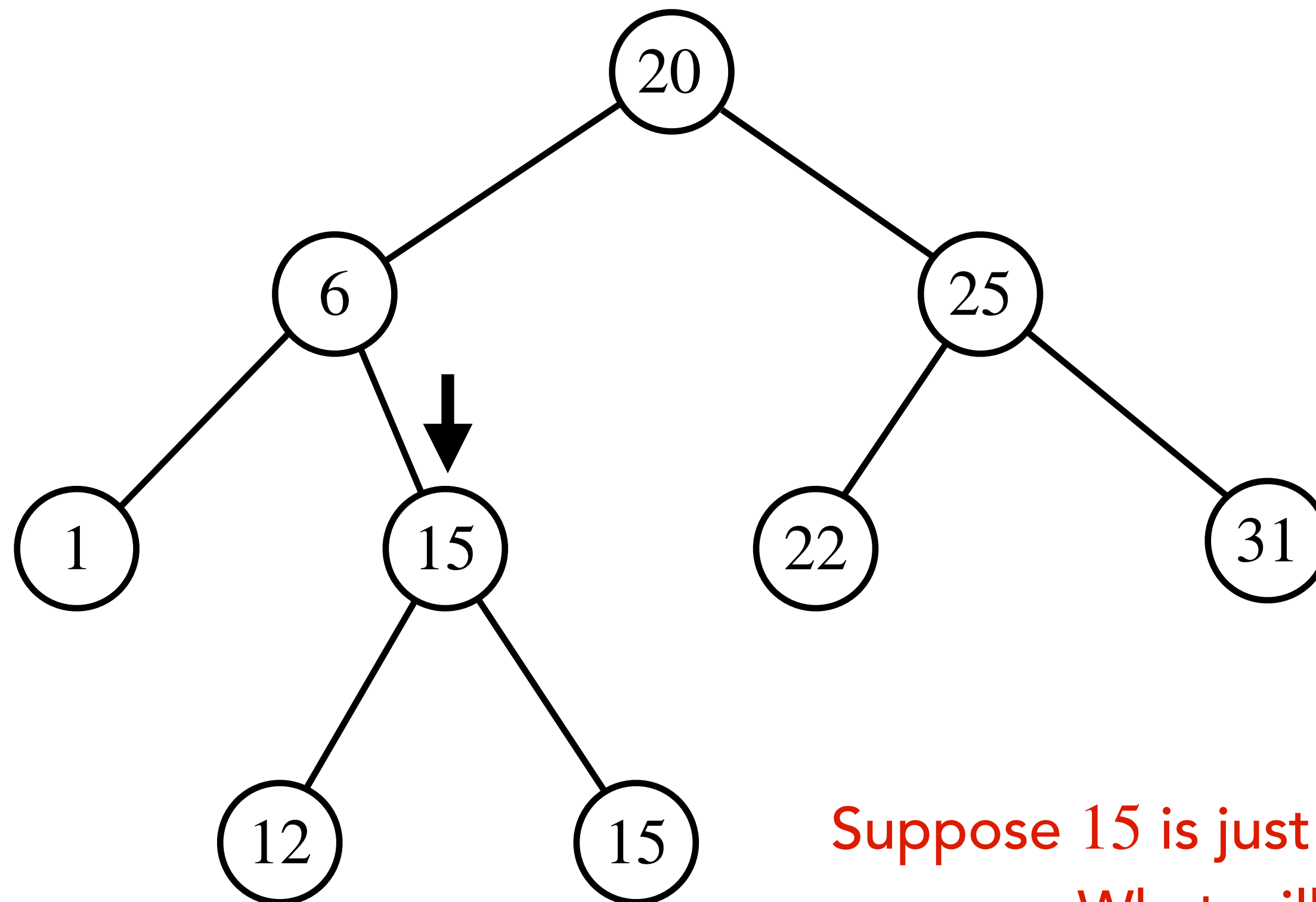
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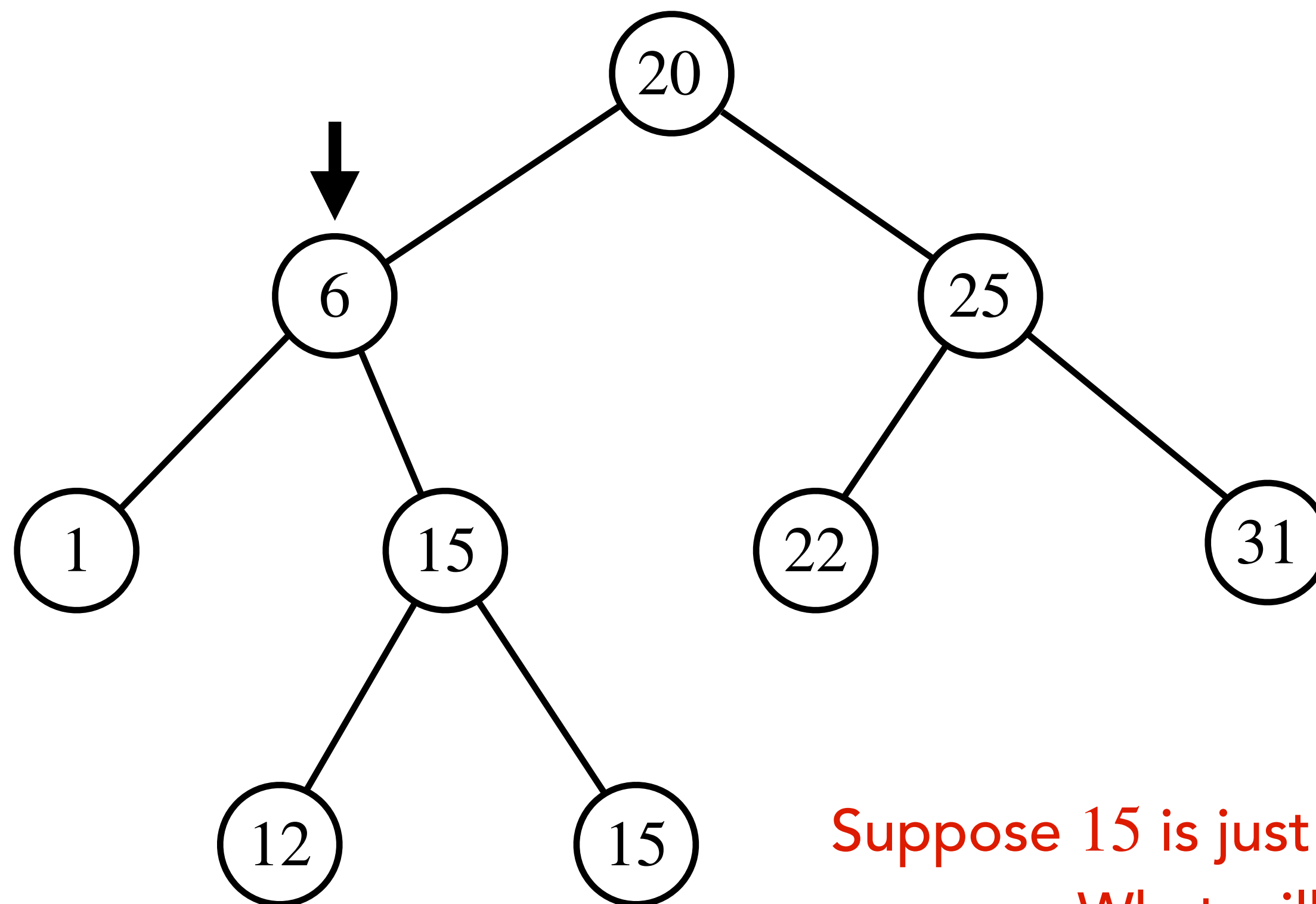
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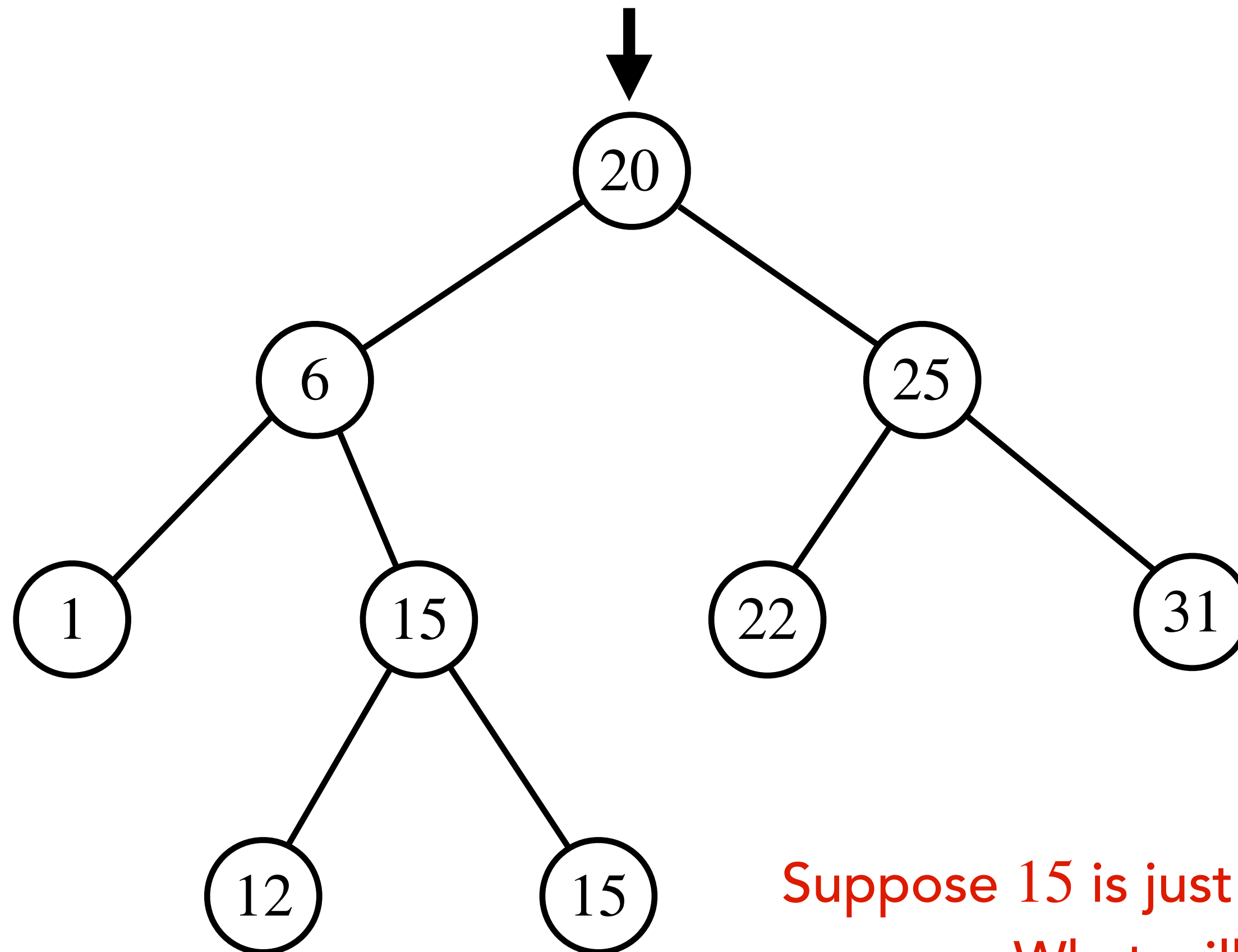
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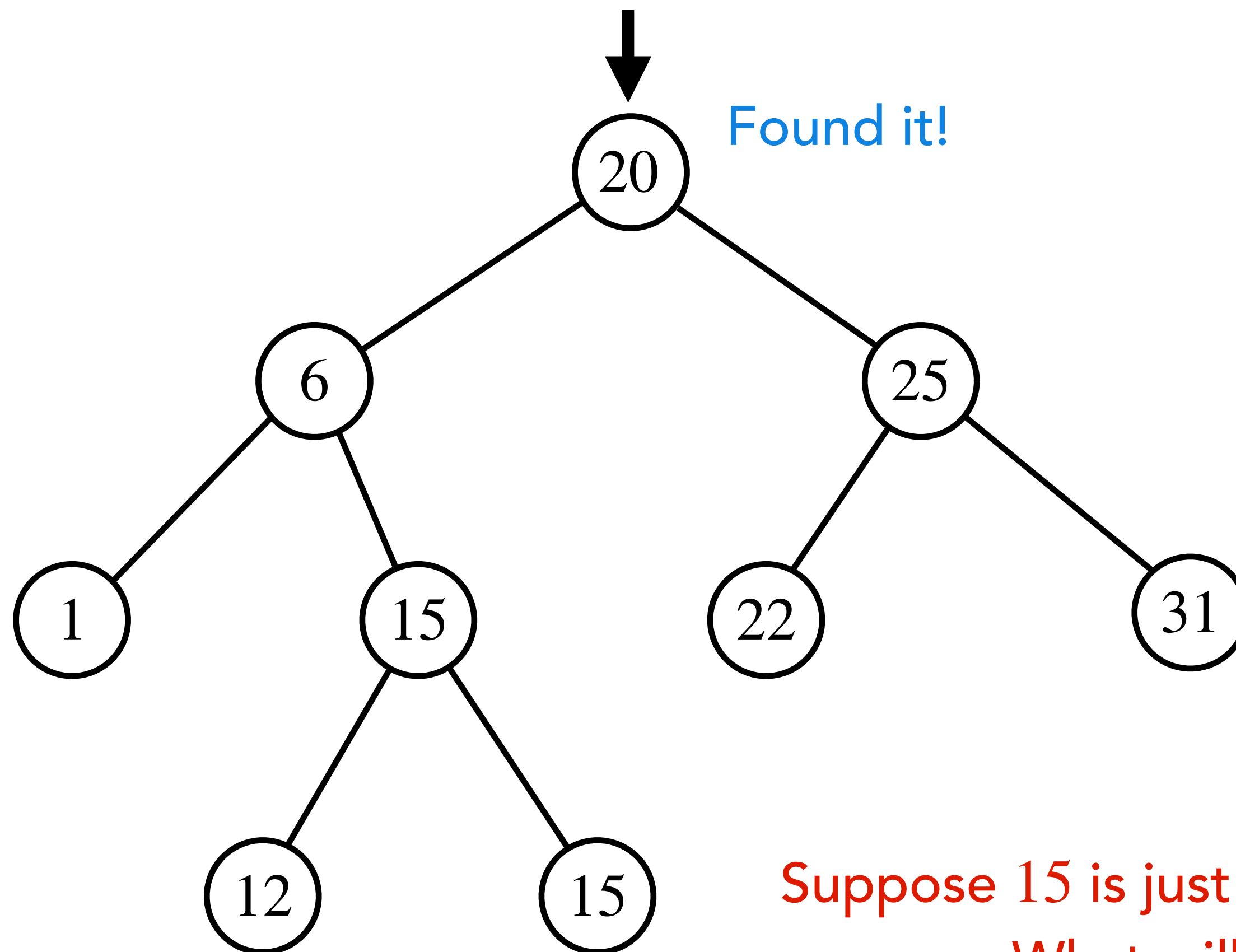
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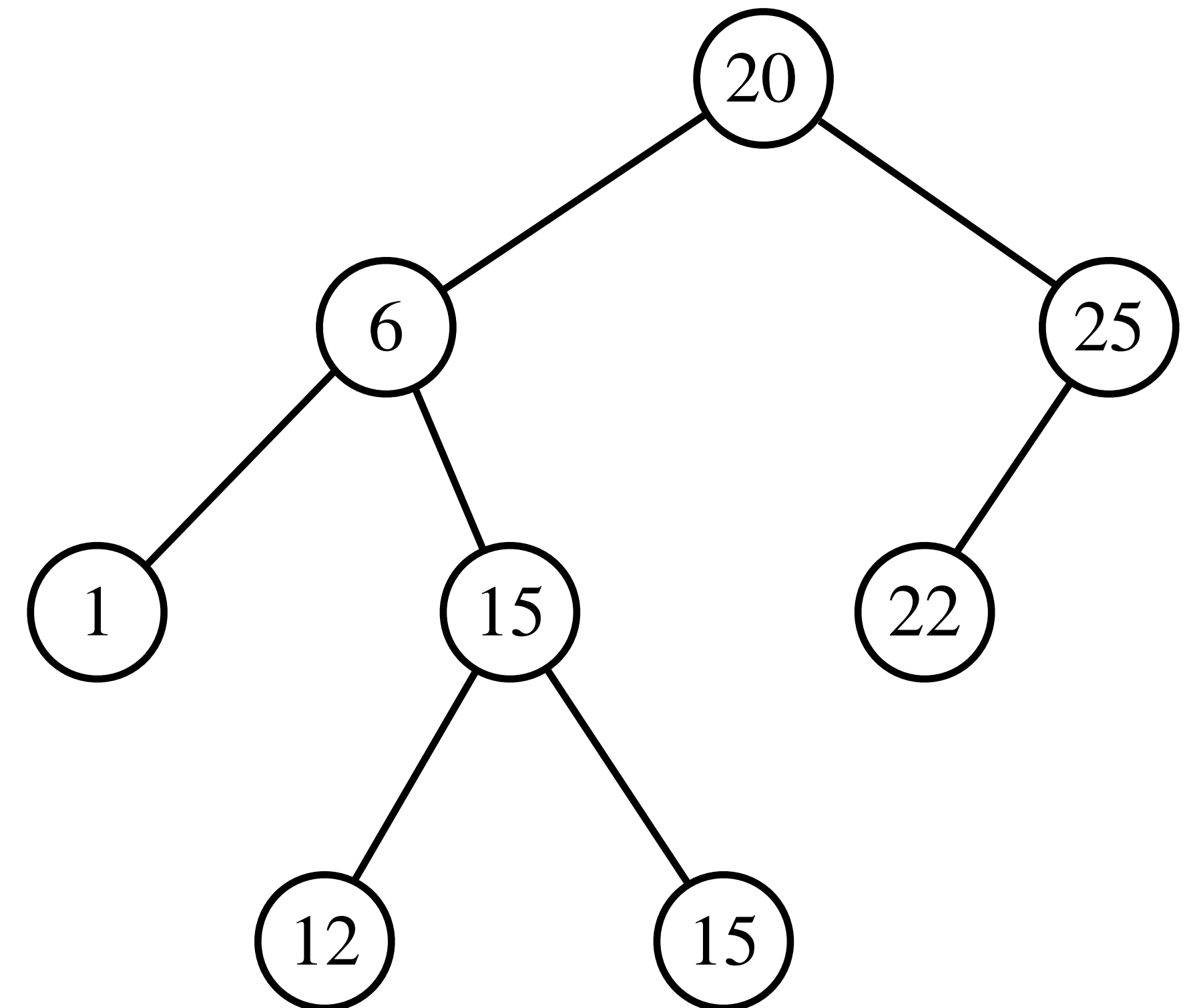
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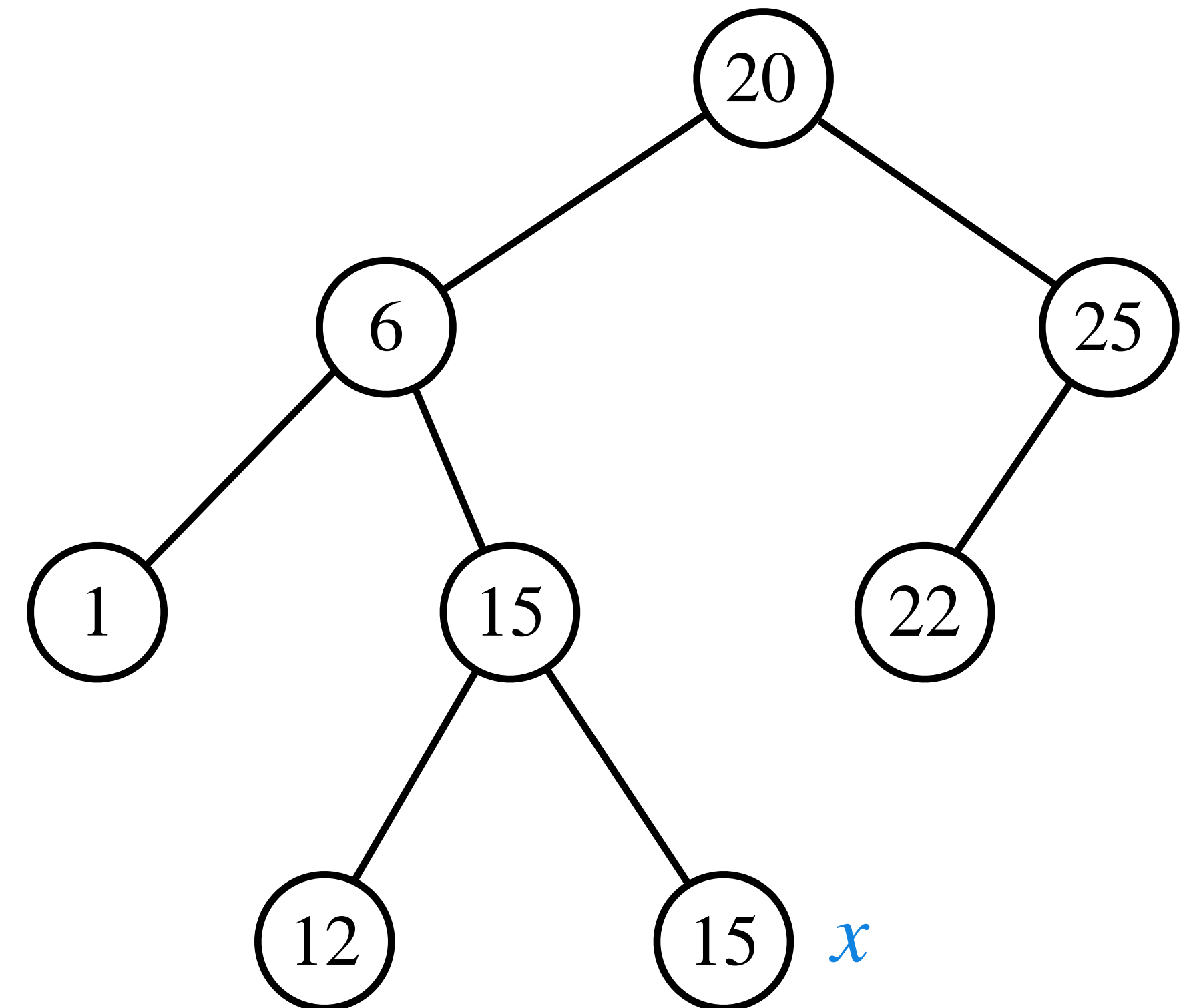
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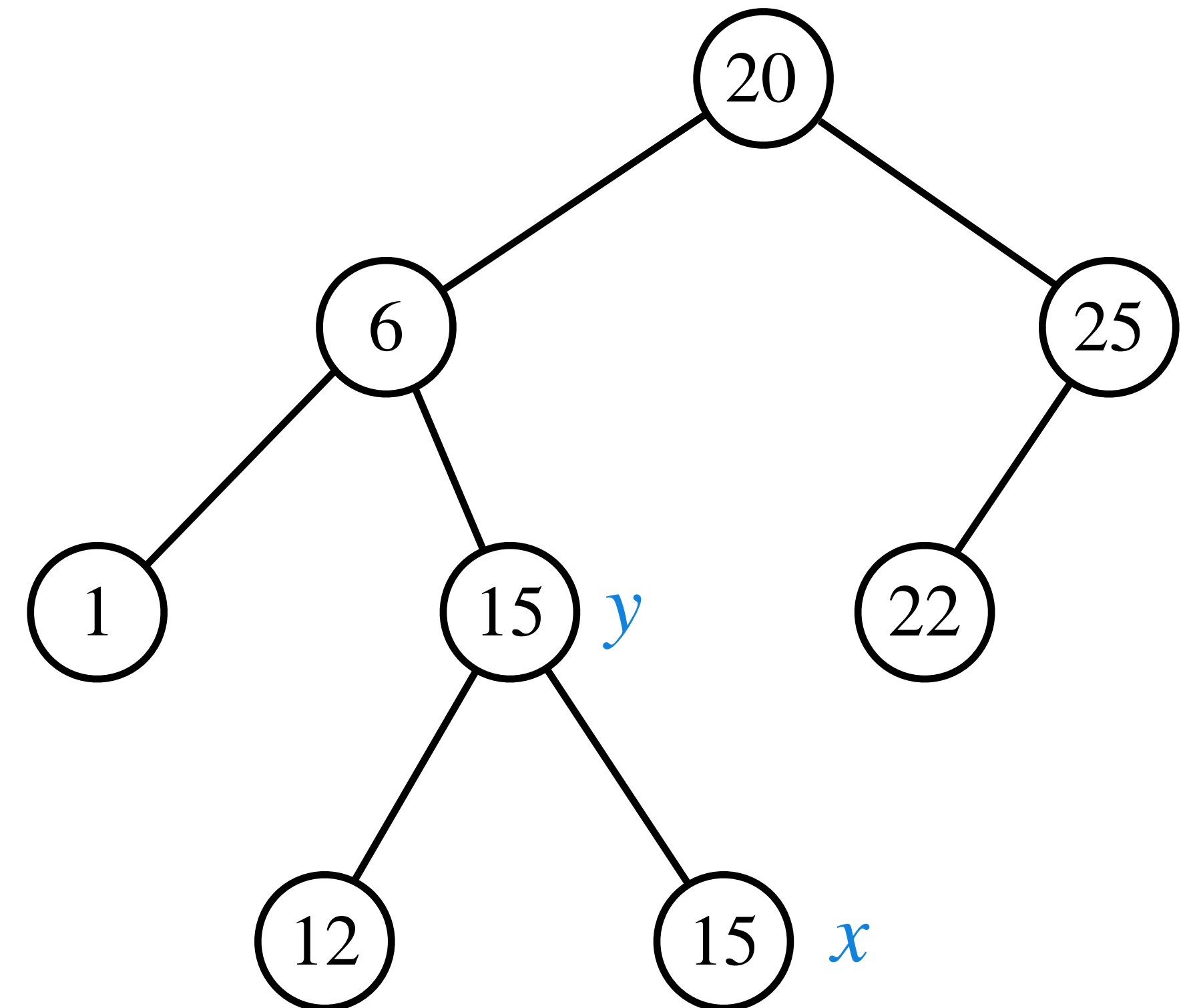
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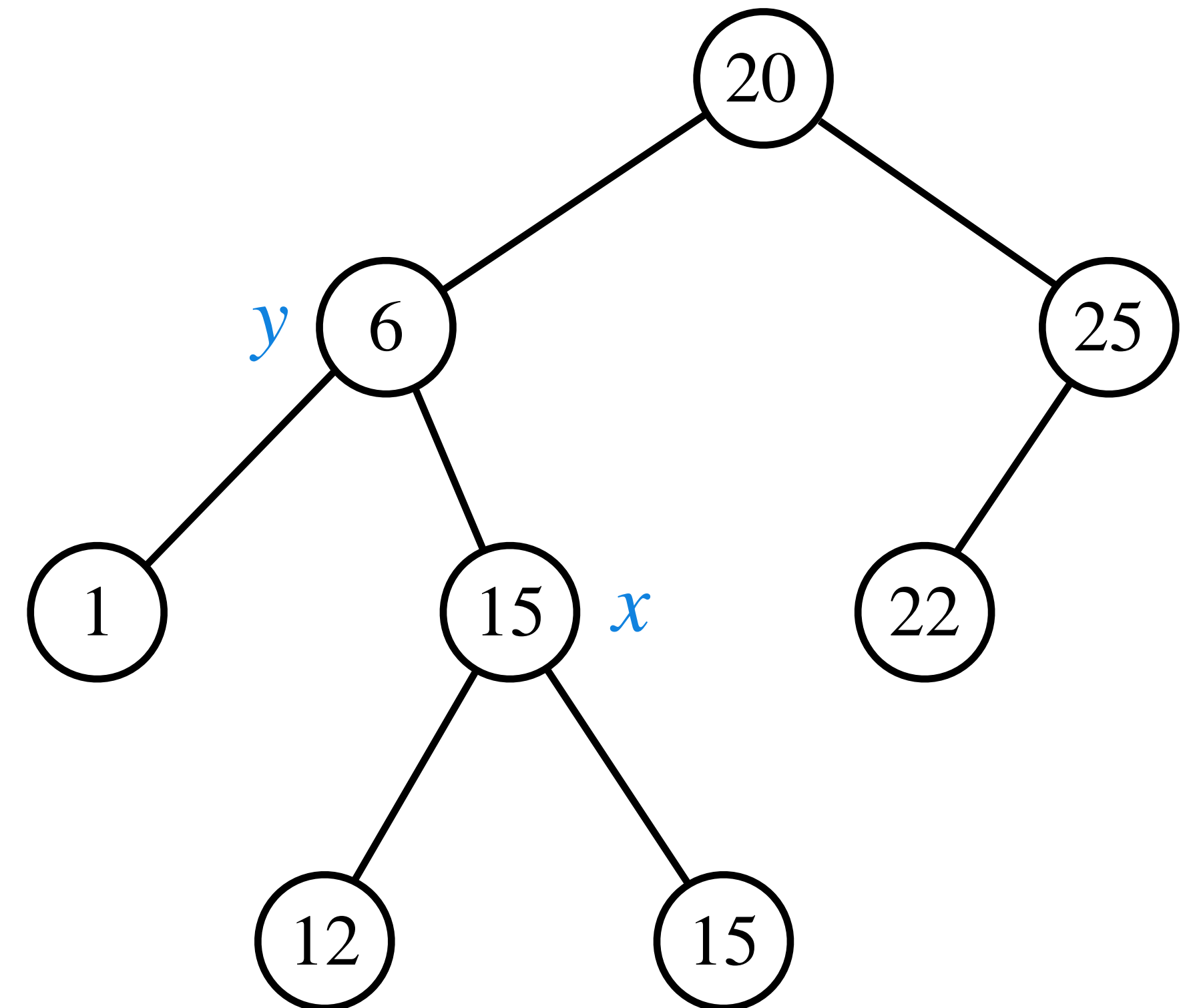
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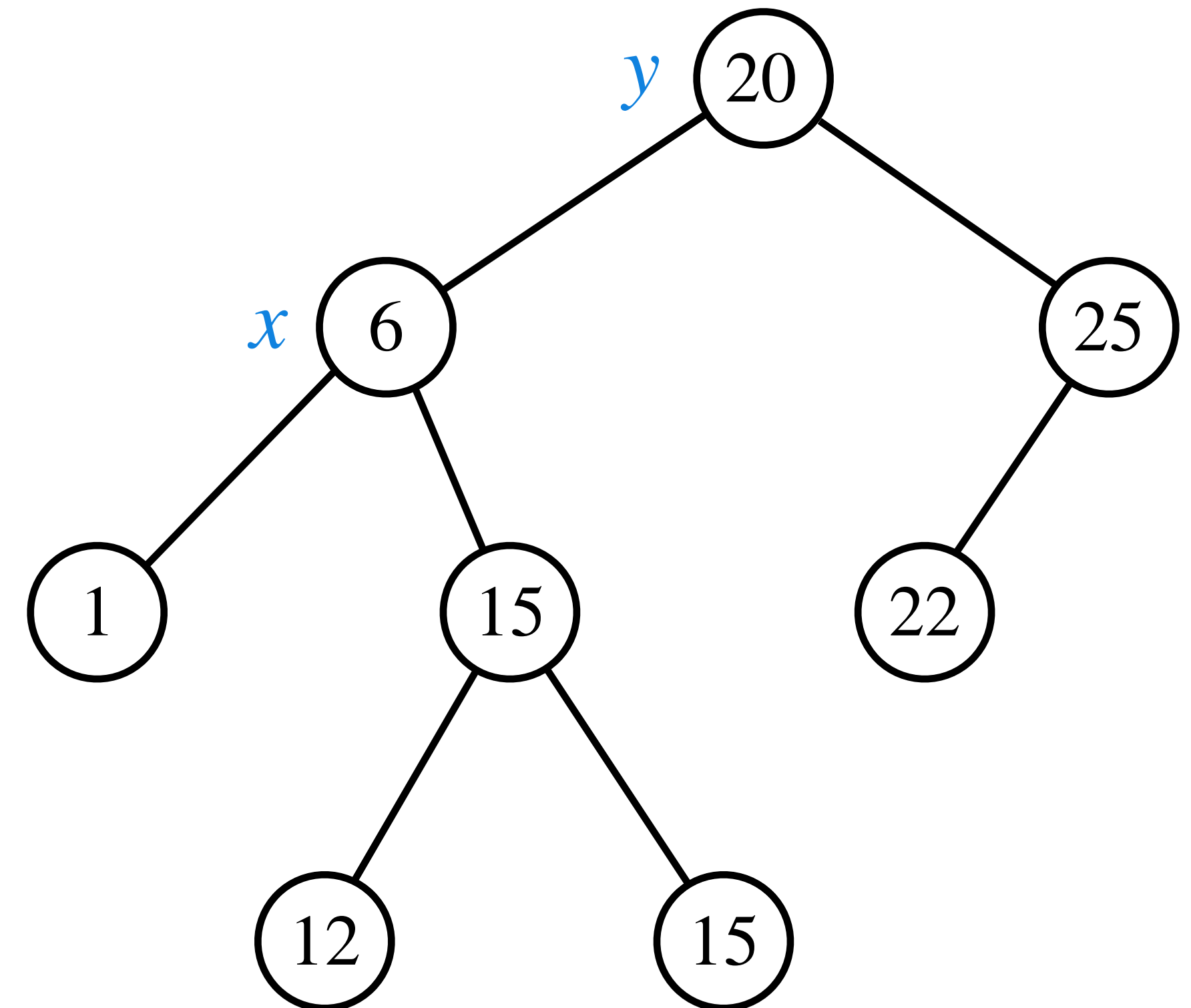
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Runtime: $\Theta(h)$, where h = height of T .

